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FORMAL METHODS FRAMEWORK

WetStone Technologies

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1 Period of Performance

This report reflects performance from 5/26/99 through 10/26/99.

2 Detailed Program Schedule

The following represents the schedule for this project:

Progress		Months				
	Jun	Jul	Aug	Sep	Oct	Nov
Perform initial research & project setup	4					
Form Collaborations / Research Tools / Distribute questionnaires		j.				
Examine tools for possible framework properties			e feeten			
Identify framework properties					沙漠	
Prepare Final Report						

3 Background

Our survey of the current practices in formal methods in academia and industry [Barj98] indicates that formal methods (FM) are a promising technology that is eliciting more and more industrial interest. Major issues in software and hardware industry are complexity and size, and current practices such as simulation cannot perform to the desired level of satisfaction anymore.

In the hardware industry, formal tools are popular and adopted in standard engineering practice. Many tool vendors such as Crysallis or Synopsis make formal tools and/or integrate formal tools in their commercial CAD toolkits. For Example, Cadence is currently producing a "Verification Cockpit" toolset. Incentives are: high cost of design errors, standard notation (VHDL/Verilog), and use of standard tools. Formal methods replace simulation, with the prevalent use of model checking to reveal errors.

In the commercial software industry, there is none or very little use of formal methods. The barriers include: product patches are distributed electronically, software is written in many languages, there is very little use of any tools, and software engineering is not a discipline based on formalism and mathematics such is digital design.

High assurance and telecommunications software industry use formal tools to some degree, with their use increasing. Telecommunications industry is driven by (often international) standards

compliance and need for test—case derivation. Information security industries, such as electronic commerce and banking, network security, and military applications are motivated by the virtue of information as commodity, with tangible material and strategic cost. Safety critical applications applications, such as avionics, medicine, railroads, and nuclear power applications are motivated by having human lives at stake.

Most formal methods practitioners agree that many additional steps are needed to take formal methods from research to industrial practice. Most commonly mentioned features include: infrastructure, such as robust and supported tools, easy to use, with verified libraries; publicizing success stories; and user education.

Some preliminary work can be done in order to make formal methods more approachable to users. IEEE Formal Methods Planning Group met in an open meeting in November 1998 at SRI International, Menlo Park, CA, to discuss what steps, if any, can be taken towards standardization of formal methods. The consensus was that standardization is premature, and that it would be necessary to collect information on the existing formal tools and somehow classify them, and collect and standardize formal methods terminology. The project described in this report addresses those concerns.

We attended World Congress on Formal Methods, which was attended by about 500 formal methods specialists from all over the world. During 1.5 hr meeting called IEEE Formal Methods Planning Group Birds-of-a feather meeting, formal methods experts discussed what needs to be done in formal methods, using our work on tool classification and taxonomy. The resulting recommendations are included in the conclusion.

A long term goal of WetStone Technologies, Inc. is to produce a robust, industrially usable Formal Methods Framework (FMF) that is populated by several formal methods and tools. This framework must be extensible, scaleable, and general enough to address a range of application problems but specific enough to address desired application domains. An undertaking of this size would require partnership between several teams with different expertise and several years of work. In this effort, we are taking the first step by outlining the preliminary work necessary to pave the way for the creation of the fully developed Formal Methods Framework.

3.1 Existing Tool Classification and Terminology Documents

Some documents and databases which outline formal methods terminology, tools used and experience reports already exist. We will not discuss databases of links to various tool pages, such as [BoweWWW], but rather databases which attempted to classify tools based on some predetermined criteria.

Formal Methods Europe (FME) is "an organization supported by the Commission of the European Union, with the mission of promoting and supporting the industrial use of formal methods for computer systems development." FME organizes seminars and a yearly international symposium, and produces a newsletter. FME's web page [FME] contains some case studies, formal methods database, and a tools database. The case studies database seems not to be up to date, and the tools database seems not to be up to date, with the latest additions in 1997, although the web page claims 6-month updates. The tools database contains about 60 international tools and is, in our opinion, suitable for a quick overview of tools. The tools are classified by the following categories:

- Tool name
- Usage and applicability
- Languages supported
- List of applications (if available)
- Functionality: yes/no answers to the following:
 - Syntax checking
 - Static semantics
 - Animation.execution
 - GUI
 - Pretty-print
 - Typechecking
 - Proof support
 - Refinement
 - Test-case generation
- Environment, number of installations, last update
- Contact
- Availability
- Description.

European Workshop on Industrial Computer Systems (EWICS) Formal Methods subgroup produced documents that contain some formal methods terminology, formal methods database, and a classification of methods by their theoretical basis [EWICS98]. EWICS formal methods database is relatively current (dated June 1998) but it contains only CCS, COLD, OBJ, SAGA-LUSTRE, Z, RAISE, B and VSE formal methods, and it focuses more on methods than on tools. The methods are classified as:

- Formal method name
- Summary
- Applications
- Properties
- Relation to other formal methods
- Theoretical basis
- Tools
- Appraisal:

- Maturity
- Availability
- Strength
- Industrial experience
- Tool availability
- Application and experience matrix
- Tools matrix
- Bibliography

Craigen, Gerhart and Ralston have published "An International Survey of Industrial Applications of Formal Methods" [CrGeRa93], which contains much valuable information but is dated as 1993. [ClWi96] paper has some experience reports as well and is more recent. Experience reports need to be kept in an up-to-date database available on the web.

Some definitions of formal methods terms are published in the following reports and databases, such as: NASA's Formal Methods Guidebooks [NASA97, NASA98]; EWICS' Guide [EWICS98]; Laprie's report "Dependability: Basic Concepts and Terminology" [Lapr]; "Dictionary of Algorithms, Data Structures and Problems" [Black99], compiled by Paul Black for CRC Dictionary of Computer Science, Engineering and Technology; and Rushby's technical report on "Formal Methods and Their Role in the Certification of Critical Systems" [Rush93, Rush95]. Various formal methods terminology is scattered throughout the published literature, such as [ClWi96]. Effort is needed to collect the terminology as is used today and converge it into a common terminology, i.e. formal methods "lingua franca."

4 Project Activity

The immediate goals for this project were to:

- 1. Collect terminology and develop a taxonomy of terms used in formal methods
- 2. Classify a subset of formal tools.

Our intent is to contribute to a more widespread use of formal methods by making formal methods more accessible and understandable to potential users, including industrially-oriented users new to formal methods. We developed a questionnaire that should help potential and new users assess what tools are available for their needs. The questionnaire was used to collect information on selected tools, and develop classification and taxonomy based on the collected information.

We have presented this work at World Congress on Formal Methods (FM'99) during IEEE Formal Methods Planning Group Birds-of-a-Feather meeting. Discussion ensued that points in the direction of future work and confirms the orientation towards industrial practitioners. The main points of the discussion are outlined in the "Conclusions" section.

4.1 Classification of Formal Tools

We compiled information on the best-known and widely used formal tools, with emphasis on tools aimed at industrial practitioners without extensive formal methods expertise. The tools are:

- 1. Theorem provers: PVS, ACL2, HOL, Larch LP tool, Z/EVES;
- 2. Model checkers: SMV, SPIN, Murphi, Concurrency Factory;
- 3. Other tools: NRL Protocol Analyzer, SCR*, Tatami.

We have devised a questionnaire to aid in collecting and classifying this information. There are many criteria for classifying the tools, based on the intended use of the tool survey. Possible audiences include tool developers, industry/users, and academia/researchers. We have assumed that the tool survey will be used by users new to the tools to aid them in selecting appropriate tools. We envisioned users who are interested in practical application of the tools and possibly do not have extensive background in formal methods. We chose the main categories for classifying the tools to be:

- 1. general description of the tool;
- 2. tool implementation (such as what language the tool is implemented in, is the tool extensible);
- 3. tool features and utilities (such as validated libraries, GUI, typechecking, prettyprinting, editing);
- 4. tool input and output;
- 5. tool applications (such as application domains, levels of abstraction;
- 6. resources required to run the tool (such as licensing, platform, operating system);
- 7. resources available (such as manuals, courses, contacts);
- 8. more specific detailed questions pertaining specifically to model checkers and theorem provers; and
- 9. open-ended questions for quick assessment of tools' strengths and weaknesses, and a list of case studies and experience reports.

We have designed the basic questionnaire and revised it based on the feedback from the Engineering Consortium, various verification mailing lists, and SRI CSL. We also modified the questionnaire for on-line filling.

For each tool, we filled the questionnaire as a "new" user, i.e. we have studied the readily available literature about the tool as if we are evaluating it for potential use. Questionnaires were then distributed to tool makers and user mailing lists for feedback. Returned questionnaires were edited for consistency between various responses. All questionnaires came back with feedback except for Larch LP, Tatami, and NRL Protocol Analyzer. (According to Jeannette Wing at FM'99, use of Larch language and tools is on the sharp decline and that might explain lack of interest in participating in this survey.)

The questionnaires are in the Appendices. We posted the questionnaires on WetStone's web page, as http://www.wetstonetech.com/fm_quest.html, and requested that it be linked to various formal methods web repositories, such as Engineering Consortium page and World Wide Web Virtual Library on Formal Methods. The questionnaires were presented at the World Congress on Formal Methods (FM'99).

4.2 Taxonomy of Formal Terms

We have examined [NASA97, NASA98] FM guidebooks, various papers on formal methods including [ClWi96] and many others, various technical reports such as [Rushby95] and combined existing definitions into a formal methods terminology. The taxonomy is application-domain independent. It is intended to satisfy a wide range of users, including practicing engineers who might not be fully trained in mathematical logic. For a more theoretical treatment of technical details involved in formal methods, a reader is referred to textbooks on logic and theoretical studies of languages such as [EWICS98]. The taxonomy is presented in the Appendix, and posted on WetStone's web page at http://www.wetstonetech.com/fm_quest.html.

5 Conclusions and Future Work

The following work is needed to move formal methods into a more mainstream practice:

- 1. A common terminology. Various differing definitions need to be converged into a common terminology to be accepted as the "lingua franca" of formal methods, and potentially standardized.
- 2. Common APIs and exchange formats for tool interoperability, potentially to be standardized.
- 3. Classification of formal methods, based on their language, method, and tool, as well as the relationship between them. Ideally, also include classification based on application domains.
- 4. Guidelines for using formal methods in industrial practice, including the following:
 - a. Overview of the state-of-the-art in formal methods practice.
 - b. A classification of tools, containing short overview and description of each tool, time-stamped and indicating if the tool is industrial strength or research prototype.
 - c. A questionnaire which users can use to guide them in selecting tools.
 - d. Experiences database, organized by application type and industry area; or, for each tool, what types of problems it was used for.
 - e. Examples done in each tool, using similar problems as benchmarks.
 - f. A catalogue of formal methods courses, training, books, and other educational resources.
 - g. "Method behind the method" for tools, i.e. how can each method/tool be used and/or what is its theoretical basis for implementation.
 - h. A bibliography of links to the above information, to be posted on a web site.
- 5. Developed "infrastructure," such as verified libraries and transition from research prototype tools into industrial strength tools.
- 6. Integration of tools into toolkits, and integration of tools into industrial process flow.

This project has accomplished items 3.b and 3.c, and produced the first draft of item 1. [Barj98] addressed item 3.a, but the overview needs to be updated yearly. The future work would be to address the remaining items.

5.1 Long-term Future Work: Formal Methods Framework

In order to integrate tools into a framework, items 2 and 3.g must be completed first. Our perspective and long-term goal is to identify robust tools that can be integrated together in a formal toolkit or added to existing toolkits. In order to accomplish that, we need to:

- Identify application or class of problems. Possible choices at this moment seem to be:
 - hardware/software co-design
 - information/networking security
 - system-level design.
- Identify collaborators. Tool integration can be achieved only with the assistance of tool makers. We need to identify collaborators that can bridge the gap between research and industrial practice.

Potential collaborators include: Derivation Systems (contact Dr. Bhaskar Bose); Dr. Perry Alexander (U of Cincinnati); SLDL project (contact Dave Barton, Intermetric Inc.) and the tool integration group at Ptolemy Project (contact Dr. Edward Lee, U of California at Berkeley).

Derivation Systems company is dedicated to making industrial formal methods products. Employees are Ph.D.-level trained in formal methods. Therefore, this company provides expertise in commercial applications of formal methods research. Furthermore, the company sells formal hardware tools, and recently has acquired software expertise in formal network assurance for secure Java applets.

SLDL (System-Level Design Language) project is an ongoing, industry-driven effort to develop a language and its tool support for describing systems-on-silicon. SLDL is intended for use by electrical engineers designing microsystems with embedded software. SLDL is of interest to our project because of its plug-and-play architecture. SLDL framework will include bridging the semantics of several existing domain-specific system languages (e.g. Esterel, SDL, and C++).

Dr. Perry Alexander has been involved in several projects that bridge the gaps between formal methods research and industrial practice, hardware and software. For example, CEENS project (sponsored by Air Force), SLDL project, and HEPE project (sponsored by DARPA/ITO). CEENS project had the goal to develop methodology and tools necessary to support board and module level of electronic integration and develop a commercial products. The project involved collaboration between Dr. Alexander and commercial companies TRW, Motorola, and Mentor Graphics, and included industry review board. HEPE project deals with high assurance heterogeneous network assurance prediction, and thus provides with software experience.

Some of the tools we have examined already integrate with other tools, for example PVS integrates with SCR*, which integrates with SPIN. The trend is between integration between theorem provers and model checkers, such as in PVS and SMV.

There are many ways to integrate tools. Ideal toolkit would consist of a "stack" of tools that can address various levels of abstraction to aid in development by transformation. Ideally, tools would be able to share common data. In practice, tools have been integrated based on shared APIs and sockets (such as in Z/EVES); or common meta-language (such as in UniForm and Express IT toolkits, and many commercial non-formal CAD toolkits) or some logic as the logical framework (such in Maude). We envision a formal methods framework that integrates several tools in an open environment. The tools should be either extensible or with provided API, and contain many "convenience" non-formal tools, such as typecheckers, editors and prettyprinters, as well as validated libraries. What is needed is more validated libraries for theorem proving, and macros of temporal logic formulas for model checking. Our goal would be to integrate theorem proving and model checking in an efficient way.

Tools which look promising for such integration are PVS, SCR* and SPIN, since they have already begun their integration. For example, PVS already contains a model checker, but it would be more efficient from a user's point of view to not have to learn another tool, e.g. an experienced SPIN user should be able to supply input to PVS and vice versa.

An outcome of this work would be to produce guidelines on how to express various properties in various tools, which is one of the features states as "needed" by the formal methods community.

An advantage of combining tools such as PVS and SPIN is the possible extent of cooperation from tool vendors and users. PVS distribution does not contain source code, so it is not user extensible, but PVS is a product of a commercial company and there are human resources available to extend the tool, even though PVS itself is free. SPIN source code is freely available and often extended by users. The combination of the two could result in a tool suite that is free and capable of integrating model checking and theorem proving in an effective way.

5.2 Travel

Date	Destination	Purpose of Trip

5.3 References

[ACL2] Applicative Common Lisp (ACL2) home page, http://www.cs.utexas.edu/users/moore/acl2/index.html.

- [Barj98] Milica Barjaktarovic, "The State-of-the-art in Formal Methods," AFOSR Summer Research technical report for Rome Research Site, AFRL/IFGB. http://www.wetstonetech.com/fm_quest.html.
- [Blac99] Paul Black, ed., "Dictionary of Algorithms, Data Structures, and Problems," compiled originally for the CRC Dictionary of Computer Science, Engineering and Technology. http://hissa.ncsl.nist.gov/~black/DADS.
- [BoweWWW] Jonathan Bowen (webmaster): "WWW Virtual Library on Formal Methods", http://www.comlab.ox.ac.uk/archive/formal-methods/, links to individual tools' pages.
- [ClWi96] Edmund M. Clarke, Jeannette M. Wing, at. al. "Formal Methods: State of the Art and Future Directions." ACM Computing Surveys, 28(4):626-643.
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- [EWICS98'] European Workshop on Industrial Computer Systems, Technical Committee 7 (Safety, Reliability, Security), Formal Methods subgroup, "Guidance on the Use of Formal Methods in the Development of High Integrity Industrial Computer Systems." Part III, "A Directory of Formal Methods," working paper 4002, June 1998. http://www.ewics.org.
- [Factory] Concurrency Factory home page, http://www.cs.sunysb.edu/~concurr
- [FME] Formal Methods Europe home page, http://www.fme-nl.org. or http://www.fme-nl.org.
- [HOL90] HOL home pages, http://www.comlab.ox.ac.uk/archive/formal-methods/hol.html
- [Lapr] Jean-Claude Laprie, "Dependability: Basic Concepts and Terminology." Laboratory for Analysis and Architecture of Systems (LAAS) CNRS, LAAS report No92043. http://www.laas.fr.
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- [NASA98] "Formal Methods Specification and Verification Guidebook for Software and Computer Systems, Volume I: Planning and Technology Insertion" [NASA/TP-98-208193], 1998. http://eis.jpl.nasa.gov/quality/Formal_Methods/

[NASA97] "Formal Methods Specification and Analysis Guidebook for the Verification of Software and Computer Systems, Volume II: A Practitioner's Companion"

[NASA-GB-001-97], 1997.

http://eis.jpl.nasa.gov/quality/Formal Methods/

[NRL] NRL home page, http://www.itd.nrl.navy.mil/ITD/5540/projects/crypto.html

[PaulWWW] Larry Paulson (webmaster): http://www.cl.cam.ac.uk/users/lcp/hotlist#Systems

[TalcWWW] Carolyn Talcott (webmaster):

http://www_formalstanford.edu/clt/ARS/ars-db.html

[PVS] PVS home page, http://pvs.csl.sri.com

[Rush93] John Rushby, "Formal Methods and Their Role in the Certification of Critical

Systems", SRI International Technical Report CSL-93-7, March 1993.

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[Rush95] John Rushby, "Formal Methods and Their Role in the Certification of Critical

Systems", SRI International Technical Report CSL-95-1, March 1995.

http://csl.sri.com/csl-95-1.html.

[SCR] SCR home page, http://www.chacs.itd.nrl.navy.mil/SCR

[SMV] SMV home page, http://www-cad.eecs.berkeley.edu/~kenmcmil

[Spin] Spin home page, http://netlib.bell-labs.com/netlib/spin/whatispin.html

[Z/EVES] Z/EVES home page, http://www.ora.on.ca/z-eves.

Appendix A – Questionnaires

ACL2

************************* ACL2 ***********************

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: model checker theorem prover mechanized proof assistant _X other: integrated toolkit: logic, mechanized proof assistant,
executable model environment.
o Application domain(s) or class(es) of problems originally intended. Formal verification o digital systems.
Building executable models that can be run and/or symbolically
executed.
o Intended audience. Engineers and mathematicians working on industrial-strength
applications. More generally, anyone wanting to reason about formal models.
o Language(s) and/or technique(s) that the tool is based on. ACL2 logic (a subset of first-order applicative Common Lisp, i.e.
excluding non-applicative aspects such as higher-order functions, circular structures, and Common Lisp Object System).
o Reasoning mechanisms used for the tool.
Mathematical induction, rewriting, decision procedures (equality, BDDs, linear arithmetic), heuristics
o Comparable languages/tools.
<pre>HOL, PVS, (Pc-)Nqthm. ACL2 is industrial-strength successor of Boyer-Moore theorem prover Nqthm).</pre>
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation. Applicative Common Lisp (Allegro, GCL, Lispworks, Lucid, MCL).
o How extensible and/or customizable is the tool. X source code given
_X source code given _X tool implemented in a public-domain language
X other: users post libraries
Features enabling modification include
extensive comments in sources and applicative
coding style (e.g. no global variables)

3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): GUI _X Library of standard types, functions, and other constructions _X the library is validated
The extent of the library is (speaking from the point of view o a potential user): not very comprehensive _X_ reasonably comprehensive quite comprehensive
X Editing and document preparation tools GNU EmacsACL2 event files can be published in LaTex, HTML, Scribe, or ASCII text. Formatting is user-extensible Cross-referencing Browsing
Requirements tracing X Incremental development across multiple sessions Change control and version management X Consistency checking (via the "encapsulate" form)
_X Completeness checking
o How interactive/mechanized/automated is the tool. _X fully automated
4. TOOL INPUT AND OUTPUT
o Tool supports these models: _X synchronous _X asynchronous X mixed
o Input to the tool. Model in ACL2 and proof hints. o Output from the tool. Proof results. Execution results.
o The language used for input to the tool has (check all that apply): _X formal semantics X modern programming language constructs (e.g. if-else):

<pre>strong typing X modularity X hierarchical design X parameterization (in the sense that functions can be parameterized) communication between processes</pre>
5. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply): _X requirements _X design specification _X implementation _x test derivation (not part of the system, but conveniently user-implementable) _X RTL _X netlists transistor level _X other: In principle, any level can be addressed, but some levels would require more work than others.
o Has the tool been integrated with other tools? no
<pre>6. RESOURCES o Resource requirements for the tool: UNIX versionSun OS, Linux Windows version Mac version _X Memory:at least 16MB, preferably at least 64MB</pre>
o Cost, rights and restrictions: free, no license _X free, license required

other:
o User background prerequisites (check all that apply):
_X BS degree
MS degree
Ph.D. degree
X knowledge of logic
X first-order
high order
familiarity with a high-level programming language
familiarity with process algebra
familiarity with temporal logic
X other: minimal familiarity with Common Lisp
o User's learning curve, if all prerequisites are met:
one month
two months
X less than six months
other:
months
o Tool support
_X upgrades/maintenance
Last version produced at this date: ACL2 v.2.4, 1999
X manual
X on the web
_X training
(tutorials on the web)
_X listserv
mailing list
mailing list _X_ dedicated conference(s)/workshop(s)
(One held in March 1999; next is anticipated in Oct. 2000)
human "help line"
x book(s)
(To appear in 2000).
_X journal/conference publications
_X other: _bug reports to acl2@lists.cc.utexas.edu
libraries, hypercard on the web
o Current contact.
http://www.cs.utexas.edu/users/moore/acl2/index.html
acl2@lists.cc.utexas.edu (subscribe to acl2-
request@lists.cc.utexas.edu)
requestarises.co.ucexas.cau/
7. QUESTIONS APPLYING TO MODEL CHECKERS ONLY
7. QUESTIONS AFFEIING TO MODEL CHECKERS ONLY
O Verification mechanism(s) (check all that apply):
equivalence
modal logic
modar rogre

temporal logic
system or process invariants
built-in support for checking for:
deadlock
livelock
other:
other:
o Tool supports (check all that apply):
optimization and state reduction mechanism
using
simulator:

interactive
random
feedback on in what state verification failed
trace leading to the state
8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
o Degree of proof mechanization.
fully mechanized
_X partially mechanized
o Support for developing and viewing the proof.
Prover gives output showing the progress of the proof that users
typically inspect in order to develop appropriate lemmas (rules) to
assist in subsequent attempts. An interactive loop allows finer
control of the proof process, as does a tool for monitoring the
rewriter. "Proof trees" provide a sort of outline mode for the
proof that can ease browsing.
o Presentation of proof to the user.
The proof is presented as formulas that the prover is attempting to
reduce
to "true".
o Tool supports (check all that apply):
X automated support for arithmetic reasoning
_X automated support for efficient handling of large propositional
expressions
_X automated support for rewriting
_X possible to use lemmas before they are proved.
_X possible to state and use axioms without having to prove them.
_X new definitions can be introduced and existing definitions
modified during proof
(at least, if "during proof" is interpreted as "during
the proof effort" then this is done all the time)
facilities for editing proofs
X the foundations (i.e., all axioms, definitions, assumptions,
lemmas) of the proof are identified
Caveat to the above: Some of the basic foundations are
collapsed, e.g., as "trivial observations"
_X reasonably easy to reverify a theorem after slight changes to
the specification

9. OPEN-ENDED QUESTIONS

o Strengths of this tool.

Industrial-strength tool.

Built and based on a programming language, so models can be symbolically

executed, run, and theorem-proved.

State-of-the-art heuristics and efficiency for inductive theorem proving.

o Limitations of this tool.

Reasoning directly about quantified notions can be very awkward. Learning curve.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Digital systems verification.

Bridging the gap between current practice (simulation) to the goal practice (formal verification) using symbolic execution, or less ambitiously, by providing a formal language for reasonably efficient simulation.

o Applications that the tool was used for - case studies, examples, success stories.

See http://www.cs.utexas.edu/users/moore/publications/acl2-papers.html. Examples:

industrial microprocessor AMD5K86 and K7 floating-point verification, Motorola CAP DSP design.

Verification of COBOL Year 2000 conversion rules.

References:

HOL

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: model checker theorem prover _X_ mechanized proof assistant other:
o Application domain(s) or class(es) of problems originally intended. General - from formalizing pure mathematics to verification of industrial hardware. Has been used for hardware and software verification.
o Intended audience. General.
o Language(s) and/or technique(s) that the tool is based on. Higher-order logic interfaced to Standard ML as the meta language.
o Reasoning mechanisms used for the tool. Higher order logic, using predicate calculus with terms from the typed lambda calculus (i.e. simple type theory).
o Comparable languages/tools. ACL2, Eves, Isabelle, Nqthm, LAMBDA, LP, Nuprl, PVS ProofProver (commercial implementation of HOL used fo reasoning about Z specifications)
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation. Standard ML (Moscow ML for HOL98, New Jersey ML for HOL90). A non-standard ML for HOL88.
o How extensible and/or customizable is the tool. _X source code given _X tool implemented in a public-domain language not extensible by user other:
3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): _X GUI (as a downloadable extension to HOL) _X Library of standard types, functions, and other constructions _X the library is validated

x	a potential user): not very comprehensive reasonably comprehensiveX quite comprehensive Editing and document preparation tools _emacs interface (as a downloadable extension)
	Cross-referencing Browsing Requirements tracing Incremental development across multiple sessions Change control and version management Consistency checking Completeness checking Other:
_x	ractive/mechanized/automated is the tool. fully automated user guided other:
4. TOOL INF	OUT AND OUTPUT
o Output fr Proof o The langu _X	the tool. er-order logic proof description. com the tool. goals proved or not. age used for input to the tool has (check all that apply): formal semantics modern programming language constructs (e.g if-else):
_x _x _x	strong typing modularity hierarchical design parameterization built-in model of computation other:
5. TOOL APP	LICATION
_x	on level that the tool can address (check all that apply): requirements design specification

	_X implementation	
	test derivation	
	X RTL	
	_X netlists	
	V transistor level	
	_X transistor level _X other: _mathematics	
	_X Other: _mathematics	
	(in principle, every level can be addressed, but lower levels	
	require more work)	
0	Has the tool been integrated with other tools?	
	no	
	X yes	
	with _Isabelle	
	with ProofProver	
	with _CHOL, non-specialist user interface to HOL	
	with	
	do not know	
	Note: Many extensions and interfaces, such as GUI, Emacs.	
	Many embedded languages, such as Z, CCS.	
	company contracts and an	
_	DECOUNCES	
0	. RESOURCES	
0	Resource requirements for the tool:	
	UNIX versionprecompiled binaries for Sun3, Sun4, MIPS, Alpha_	
	Windows version	
	Mac version	
	Memory:	
0	Cost, rights and restrictions:	
_	X free, no license	
	free, license required	
	for educational and research use only	
	nominal distribution charge	
	fee for underlying tool(s)	
	flat license fee	
	per user license fee	
	royalties per use	
	other:	
_	User background prerequisites (check all that apply):	
U		
	BS degree	
	X MS degree	
	_X PhD degree	
	_X knowledge of logic	
	first-order	
	X high order	
	familiarity with a high-level programming language	
	familiarity with process algebra	
	familiarity with temporal logic	
	other:	
0	User's learning curve, if all prerequisites are met:	
	one month	
	two months	
	less than six months	
	X other	

6 months	
o Tool support	
X upgrades/maintenance	
Last version produced at this date: HOL98	
_X manual	
x_ on the web	
_X training	
(courses at various locations, lectures and tutorials on the well	b)
X listserv	
X mailing list	
X conference(s)/workshop(s)	
(annual international intercontinental conference TPHOL)	
human	
X book(s)	
X journal publications	
Xother:web pages with code depositories and ftp/faq	
archive	
HOL2000 initiative, to design next generation	
HOL-like provers	
special journal issues related to HOL	
user meetings	
very extensive documentation (tutorial, description,	
manual, manual for each supported library, primer for	
beginners, notes, user manuals, applications)	
bug/problem reports: hol-supprt@cl.cam.ac.uk	
o Current contact.	
http://www.cl.cam.ac.uk/Research/HVG/HOL	
info-hol@lal.cs.byu.edu (subscribe at info-hol-request@lal.cs.buy.edu)	
7. QUESTIONS APPLYING TO MODEL CHECKERS ONLY	
o Verification mechanism(s) (check all that apply):	
equivalence	
modal logic	
temporal logic	
system or process invariants	
other:	
o Tool supports (check all that apply):	
optimization and state reduction mechanism(s)	
using	
simulator	
interactive	
random	
<pre>feedback on in what state verification failed</pre>	
trace leading to the state	
built-in support for checking for:	
deadlock	
livelock	
boolean propositions	
other:	

8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]

o Degree of proof mechanization.

_____fully mechanized
_X__ partially mechanized
o Support for developing and viewing the proof.

o Presentation of proof to the user (e.g., user input or canonical expressions,
with or without quantifiers).

o Tool supports (check all that apply):

_____ automated support for arithmetic reasoning
____ automated support for efficient handling of large propositional expressions
_X__ automated support for rewriting
____ possible to use lemmas before they are proved.
____ possible to state and use axioms without having to prove them.
_X__ new definitions can be introduced and existing definitions

9. OPEN-ENDED OUESTIONS

modified during proof

the specification

facilities for editing proofs

lemmas) of the proof are identified

o Strengths of this tool.

Powerful proof mechanism for formal verification, induction, infinite data sets. Active and large established user group. o Limitations of this tool.

Difficult to specify control sequences, takes a long time to learn. Less payoff for lower levels of abstraction.

_X__ the foundations (i.e., all axioms, definitions, assumptions,

_X__ reasonably easy to reverify a theorem after slight changes to

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Verification of problems containing extensive data path. o Applications that the tool was used for - case studies, examples, success stories.

Some are posted http://www.dcs.glasgow.ac.uk/~tfm/hol-bib.html Examples: embedding of various languages (e.g. Z, CCS, hardware languages); security; distributed systems; protocols; hardware; networking elements; compiler verification; real-time systems; reactive systems.

References:

Larch Prover (LP)

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
<pre>o Rough classification:</pre>
Note: each Larch specification contains two components: one written in a Larch Interface Language, which is designed for a specific programming language; and another written in Larch Shared language (LSL), which is independent of any programming language. Larch Interface Languages exists for C (LCL), Ada, Modula-3, VHDL, and others. LSL tool checks for syntax and type errors in LSL specifications, and can translate it into input files for LP. LCLint tool statically checks C programs, including common lint checks such as type inconsistencies, ignored return values, likely infinite loops, as well as assertions about assumptions in desired places in the C code ad errors in dynamic memory management.

o Reasoning mechanisms used for the tool.

Theorem proving, including forward and backward inference, equational term-rewriting, induction rules.

o Comparable languages/tools.

HOL, PVS.

2. TOOL IMPLEMENTATION

o underlying mechanism of the tool's implementation.
o How extensible and/or customizable is the tool source code given
<pre>tool implemented in a public-domain language not extensible by user</pre>
other:

3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): GUI
Library of standard types, functions, and other constructions the library is validated
The extent of the library is (speaking from the point of view of a potential user):
not very comprehensive reasonably comprehensive
quite comprehensive
Editing and document preparation tools
Cross-referencing
Browsing
Requirements tracing
Incremental development across multiple sessions
Change control and version management Consistency checking
Completeness checking
Other:
o How interactive/mechanized/automated is the tool. fully automated
X user guided
other:
4. TOOL INPUT AND OUTPUT
1. 1002 11.01 11.0 001101
o Tool supports these models:

	synchronous	
	asynchronous mixed	
0	Input to the tool.	
0	Output from the tool.	
0	The language used for input to the tool has (check all that apply)	:
	<pre>X formal semantics X modern programming language constructs (e.g. if-else):</pre>	
	_X strong typing	
	X modularity	
	X hierarchical design	
	X parameterization communication between processes	
	buffered	
	_X built-in model of computation	
	other:	
3	. TOOL APPLICATION	
0	Abstraction level that the tool can address (check all that apply): requirements	•
	X design specification	
	_X implementation	
	test derivation RTL	
	netlists	
	transistor level	
	other:	
0	Has the tool been integrated with other tools?	
	no X yes - please name tool and applications	
	with _LSL and LCLint, as mentioned above	
	with	
	do not know	
4.	. RESOURCES	
0	Resource requirements for the tool:	
	UNIX version _Intel Linux, SPARC SunOS4.1, Solaris 5.3 Windows version	-
	Mac version	

Memory:
o Cost, rights and restrictions:
_X free, no license
free, license required
for educational and research use only
nominal distribution charge
fee for underlying tool(s)
flat license fee
per user license fee
royalties per use
other:
o User background prerequisites (check all that apply):
BS degree
X MS degree
Ph.D. degree
_X knowledge of logic
_X first-order
high order
familiarity with a high-level programming language
familiarity with process algebra
familiarity with temporal logic
other:
- Warning January is all managements and make
o User's learning curve, if all prerequisites are met:
one month
two months
_X less than six months
other
o Tool support
X upgrades/maintenance
Last version produced at this date: _vs3.1b, 1999
X manual
X on the web
training
listserv
mailing list
X dedicated conference(s)/workshop(s)
human "help line" X book(s)
X journal/conference publications
X other:newsgroup comp.specification.larch
ftp archive
o Current contact.
http://www.sds.lcs.mit.edu/spd/larch/
-
6. QUESTIONS APPLYING TO MODEL CHECKERS ONLY
a Manifikashian mashanian (a) (abaab and a)
o Verification mechanism(s) (check all that apply):
equivalence
modal logic
temporal logic

	system or process invariants built-in support for checking for: deadlock livelock other:
	other:
	optimization and state reduction mechanism using symbolic simulator: interactive random feedback on in what state verification failed trace leading to the state
7. QUESTION	IS ABOUT THEOREM PROVERS [NASA98]
X_ o Support f o Presentat expressions	proof mechanization. fully mechanized partially mechanized for developing and viewing the proof. ion of proof to the user (e.g., user input or canonical hout quantifiers).
	orts (check all that apply): automated support for arithmetic reasoning automated support for efficient handling of large propositional expressions automated support for rewriting possible to use lemmas before they are proved. possible to state and use axioms without having to prove them. new definitions can be introduced and existing definitions modified during proof facilities for editing proofs the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified reasonably easy to reverify a theorem after slight changes to the specification
8. OPEN-END	ED QUESTIONS
Strengths	of this tool.
o Limitatio	ns of this tool.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

o Applications that the tool was used for - case studies, examples, success stories.

References:

PVS

For this particular tool, please answer the following questions based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: model checker theorem prover mechanized proof assistant _X other: Verification system consisting of a specification language
and support tools, including a mechanized proof checker integrated with a model checker, ground evaluator, and tabular specification tool. o Application domain(s) or class(es) of problems originally intended: Formalization and verification of requirements and design-level specifications of hardware and software systems.
o Intended audience: Anyone interested in formal support for conceptualization and debugging of algorithms, and of software and hardware systems. Both academic and industrial settings.
o Language(s) and/or technique(s) that the tool is based on: Classical, typed higher order logic augmented with predicate subtypes, dependent typing, abstract data types, and parameterized theories.
o Reasoning mechanisms used for the tool: Low-level decision procedures (including propositional simplification; ground procedures for equality, arithmetic, array, and datatype operations; and model checking) combined with user-definable, high-level proof strategies. Sequent Calculus notation. CTL model checking using mu-calculus.
o Comparable languages/tools: PVS provides more automation than a low-level proof checker (e.g., LCF, HOL, Nuprl, Coq), and more control than a highly automatic theorem prover (e.g., Otter, ACL2). PVS's capabilities are somewhat less generic than Isabelle's.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation: Common Lisp (preferably Franz Inc's Allegro Lisp) with CLOS extensions Emacs or XEmacs (version 19 or later), Tcl/Tk, LaTex.
o How extensible and/or customizable is the tool? source code given tool implemented in a public-domain language

	not extensible by the user
_x	other: The PVS environment, including Lisp, Emacs, X windows, and Tcl/Tk, are customizable. Tool makers accept and incorporate suggestions for extending/integrating PVS.
3. TOOL FEA	TURES AND UTILITIES
o Tool supp	orts the following (check all that apply):
_x	Library of standard types, functions, and other constructions _X the library is validated
	The extent of the library is (speaking from the point of view of a potential user): not very comprehensiveX_ reasonably comprehensive quite comprehensive
_x	Editing and document preparation tools
	X GNU or X Emacs _X_ Customized prettyprinting and typesetting (using LaTex) Cross-referencing
	Browsing Requirements tracing
_x	Incremental development across multiple sessions
	Change control and version management Consistency checking
	Completeness checking
	Other
	active/mechanized/automated is the tool?
	fully automated user guided
***	(simpler steps are automated)
_x	other:
	The user may also define application-specific strategies to automate the verification.
4. TOOL INP	UT AND OUTPUT
o Tool supp	orts these models:
	synchronous
	asynchronous mixed:
o Input to	
_	text consisting of a specification in the PVS language.
	om the tool:
Proof	results, status information, alltt and latex output,
	fication files, proof files. age used for input to the tool has (check all that apply):
_X:	formal semantics
	modern programming language constructs(e.g. if-else):
	if-else. let. where

structured datatypes (e.g., records, tuples, ennumerations)abstract data typestabular notation X_ strong typing X_ modularity hierarchical design X_ parameterization communication between processesbuffered built-in model of computation X_ other: Undecidable typechecking: to cope with this, the typechecker generates proof obligations, most of which are discharged automatically by the proverOverloading: PVS allows a liberal amount of overloadingAutomated support for judgements and coercions (conversions)Total vs partial functions: in PVS, functions represent total maps; partial functions are admitted within this framework via the predicate subtype mechanism.
5. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
X design specificationX implementationX test derivationX RTL netlists transistor level
logicStanford University)
6. RESOURCES
o Resource requirements for the tool: UNIX version:precompiled for Solaris 2 or higher (SPARC workstations),
Redhat Linux Windows version

	MAC VEISION
	Memory: 20 mb disk space, 50 mb swap space, 32 mb real memory
0	Cost, rights and restrictions:
	free, no license
	X free, license required
	for educational and research use only
	nominal distribution charge
	fee for underlying tool(s)
	flat license fee
	per user license fee
	royalties per use
	other
0	User background prerequisites (check all that apply):
	_X BS degree
	MS degree
	Ph.D. degree
	_X knowledge of logic
	_X first-order
	high order
	<pre>_X familiarity with a high-level programming language familiarity with process algebra</pre>
	familiarity with process argebra familiarity with temporal logic
	other:
0	User's learning curve, if all prerequisites are met:
	one month
	two months
	less than six months
	_X other
	6 months
0	Tool support:
	_X upgrades/maintenance
	Last version produced at this date: PVS 2.3, 1999
	_X manual
	X_{-} on the web
	_X training
	(tutorials on the web) listserv
	X mailing list
	dedicated conference(s)/workshop(s)
	human "help line"
	book(s)
	X journal/conference publications
	X other:
	bugs, problems, suggestions to pvs-bugs@csl.sri.com
	list of user suggestions and SRI's responses on the web
	archive, FAQ, libraries on the web
0	Contact:
	pvs-request@csl.sri.com

7. QUESTIONS APPLYING TO MODEL CHECKERS ONLY
<pre>o Verification mechanism(s) (check all that apply): _X_ equivalence _X_ modal logic _X_ temporal logic</pre>
(CTL and fair CTL)
_X system or process invariants
_X built-in support for checking for:
X deadlock _X_ livelock
X boolean propositions
X other:fairness
other
o Tool supports (check all that apply):
X_ optimization and state reduction mechanism
simulator
interactive random
x feedback on state in which verification failed (Counterexample generation is currently under development.)
8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
o Degree of proof mechanization: fully mechanized _X partially mechanized
(although finite state verification and the proof of many straightforward results are fully automatic. There is also a batch mode in which proofs may be easily rerun, and a facility for defining proof strategies to automate proofs.
o Support for developing and viewing the proof: Tcl/Tk interface to display proof trees and theory hierarchies. Proofs yield scripts that may be edited, attached to additional
formulas,
and rerun. Proofs may also be checkpointed, providing rapid access to parts of a proof the user wishes to examine or adjust.
o Presentation of proof to the user (e.g., user input or canonical expressions
with or without quantifiers): Proofs are presented in a sequent-style representation. PVS takes care to assure that the initial proof goal transparently reproduces the formula input by the user. Quantification is retained; implicit universal quantification in the user's specification is made explicit
o Tool supports (check all that apply): _X automated support for arithmetic reasoning _X automated support for efficient handling of large propositional

expressions

_X__ possible to use lemmas before they are proved. X possible to state and use axioms without having to prove them. _X_ new definitions can be introduced and existing definitions modified during proof X__ facilities for editing proofs _X_ the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified _X__ reasonably easy to reverify a theorem after slight changes to the specification _X__ other: __integration with CTL model checking ground evaluator (providing "run" speeds comparable to imperative programs) __proof strategies proof storage, replay, and checkpointing graphical display of proof trees, theory hierarchies, and prover commands __proof chain analysis proof and theory status reporting 9. OPEN-ENDED QUESTIONS o Strengths of this tool: Comprehensive, interactive environment for writing formal specifications and checking formal proofs, including tight integration of algorithmic and deductive proof technologies. Generic system well suited to, e.g., prototyping specialized strategies, embedding logics, and exploring strategies for integrating formal techniques, as well as to undertaking proofs of difficult algorithms and complex systems. o Limitations of this tool: PVS's capabilities complement, but do not compete with those of dedicated lightweight tools for specialized applications. Not industrial strength, but a mature research prototype. User learning curve. o Estimated possible uses of the tool (e.g., applications, classes of problems, stages of production cycle): Hardware verification, embedding logics, fault-tolerant algorithms, library development, invariant generation and abstraction, distributed algorithms, requirements specification and verification, security protocols, test generation. o Applications that the tool was used for - case studies, examples, success stories: Posted on http://pvs.csl.sri.com. Examples: Hardware: __Collins Commercial Avionics microprocessor design __Fujitsu high level design and validation of ATM switch

X automated support for rewriting

```
__NASA single pulser digital circuit
  IEEE 854 floating point standard
  SRT division
Distributed Algorithms:
  __FLASH cache coherence protocol
  __bounded retransmission protocol
   real-time controllers
Fault Tolerant Algorithms:
  __Fault-tolerant agreement and diagnosis protocols for various
    architectures and fault models
Embedding Logics:
   Duration calculus
   The B-method
  A real-time Hoare logic
Invariant Generation and Abstraction:
  __PVS has been used as a simplifier in several systems for
    the heuristic discovery of loop invariants for distributed
    protocols
Requirements:
   Space Shuttle flight software
  __Cassini spacecraft fault-protection software
```

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1. http://eis.jpl.nasa.gov/quality/Formal Methods/

Z/EVES

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: model checker theorem prover X mechanized proof assistant other: _ Z interface to EVES mechanized proof assistant o Application domain(s) or class(es) of problems originally intended. Analytical support for writers of Z specifications.
Formal methods courses. Various applications in safety- and security- domains. o Intended audience. Students, lecturers, researchers, commercial users interested in
rigorous specifications supported by rigorous analysis. o Language(s) and/or technique(s) that the tool is based on. Z, Verdi, s-Verdi. Verdi is a language based on untyped set theory. o Reasoning mechanisms used for the tool. General theorem proving, specifying and implementing programs, proving consistency between specification and implementation. Syntax and type checking, schema expansion, domain checking, pre-condition calculation, refinement, and general conjectures about a specification. EVES has a programming component and supports pre/post proofs, in addition to general mathematical modeling. o Comparable languages/tools. ProofPower, Cadiz and Zola.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation. Implemented in Lisp. o How extensible and/or customizable is the tool.

APIs are now defined for Z/EVES allowing for interchanges between tools.

X tool implemented in a public-domain language

____ source code given

X not extensible by user other:

only executables are distributed. 3. TOOL FEATURES AND UTILITIES o Tool supports the following (check all that apply): _X__ Library of standard types, functions, and other constructions _X_ the library is validated The extent of the library is (speaking from the point of view of a potential user): ___ not very comprehensive __ reasonably comprehensive _X_ quite comprehensive It contains all of the Spivey toolkit, which is the general basis for all Z specifications. X Editing and document preparation tools _Framemaker-based Z editor _Framemaker editor that has an API connection to Z/EVES. Cross-referencing X Browsing (to be completed soon) Requirements tracing X Incremental development across multiple sessions (to be completed soon) Change control and version management X Consistency checking X Completeness checking X__ Other: __syntax and type checking__ _schema expansion __precondition calculation_ __domain checking _proving consistncy between specification and implementation support for the Mathematical Toolkit as described in Spivey's 2nd edition of "The Z Notation" o How interactive/mechanized/automated is the tool. ___ fully automated X user guided some prover steps are automated ____ other: __ 4. TOOL INPUT AND OUTPUT o Tool supports these models: ____ synchronous asynchronous mixed o Input to the tool.

Plans are to augment Z/EVES with 3rd party developments. Currently,

z, verdi or s-verdi specification.
o Output from the tool.
Proof results. o The language used for input to the tool has (check all that apply): Note: the following paragraph refers to Verdi language: _X formal semantics _X modern programming language constructs (e.g. if-else):
<pre>X strong typing X modularity X hierarchical design X parameterization X communication between processes</pre>
3. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply): _X requirements _X design specification _X_ implementation test derivation RTL netlists transistor level _X other:mathematics
was the tool have intermed with ather tools
o Has the tool been integrated with other tools?
do not know
4. RESOURCES
o Resource requirements for the tool: UNIX versionSunOS 4.x, Linux ELF Windows version3.1,95,98,NT Mac version Memory:at least 32Mb
o Cost, rights and restrictions:

free, no license
X free, license required
for educational and research use only
nominal distribution charge
fee for underlying tool(s) flat license fee
flat license fee
per user license fee
royalties per use
other:
o User background prerequisites (check all that apply):
_X BS degree
MS degree
Ph.D. degree X knowledge of logic
X knowledge of logic
_X first-order
high order
familiarity with a high-level programming language familiarity with process algebra familiarity with temporal logic
familiarity with process argument
ramiliarity with temporal 10gic
X other: The above checked fields refer to performing proofs.
Type checking, schema expansion, pre-condition calculation,
domain checking without proof, require no knowledge of
logic
o User's learning curve, if all prerequisites are met:
one month
two months less than six months
more than six months
months
Note: Depends upon application. Type checking, schema
expansion, pre-condition calculation, and domain checking (without proof
should only take a day or two to learn. Learning to preform more serious
proofs could take several months.
o Tool support
X upgrades/maintenance
Last version produced at this date: vs.3x, due November
1999
X manual
X on the web
_X training
Course is provided.
_X listserv
<pre>mailing list</pre>
_X human
ORA will provide consulting.
book(s)
X journal/conference publications
other:
o Current contact.
http://www.ora.on.ca/z-eves/

zeves@ora.on.ca (subscribe at zeves-request@ora.on.ca)

6. QUESTIONS APPLYING TO MODEL CHECKERS ONLY	
<pre>o Verification mechanism(s) (check all that apply):</pre>	
o Tool supports (check all that apply): optimization and state reduction mechanism using symbolic simulator: interactive random feedback on in what state verification failed trace leading to the state	
7. QUESTIONS ABOUT THEOREM PROVERS [NASA98]	
<pre>o Degree of proof mechanization.</pre>	
expressions _X automated support for rewriting _X possible to use lemmas before they are proved. _X possible to state and use axioms without having to prove them. _X new definitions can be introduced and existing definitions modified during proof Would have to restart the proof. _X facilities for editing proofs _X the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified _X reasonably easy to reverify a theorem after slight changes to the specification	
8. OPEN-ENDED QUESTIONS	

o Strengths of this tool.

Rigorously developed SPARC Verdi compiler for EVES/Verdi. Synergy of an expressive writable notation (Z) with an automated Analytical engine. Useful for the Z community. o Limitations of this tool.

Limited to Z community, can take long time to learn. o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Education, safety, security.
o Applications that the tool was used for - case studies, examples, success stories.

Some are posted on http://www.ora.on.ca/biblio-welcome.html. Analysis of authentication protocols, including X.509. Design of a prototype High Assurance One-Way Link. Many proprietary applications.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1. http://eis.jpl.nasa.gov/quality/Formal_Methods/

Concurrency Factory	

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.	
1. GENERAL DESCRIPTION OF THE TOOL	
o Rough classification: model checker theorem prover mechanized proof assistant _X_ other: _integrated toolset: model checker, simulators, graphical and textual user interface, code generator	
o Application domain(s) or class(es) of problems originally intended. Concurrent systems, such as protocols or control systems; networks of finite-state processes. Industrial problems, e.g. in telecommunications industry.	
o Intended audience.	
Protocol engineers and software developers. o Language(s) and/or technique(s) that the tool is based on. GCCS, a graphical variant of the process algebra CCS aimed at specifying hierarchical networks of processes. VPL, a textual language for hierarchical networks of processes, with support for complex data and control structures. o Reasoning mechanisms used for the tool. Computing set of transitions possible for a system in a given state	
using formal operational semantics. GCCS interpreted by all the tools in the toolkit. o Comparable languages/tools. CWB, Spin.	
2. TOOL IMPLEMENTATION	
o Underlying mechanism of the tool's implementation. C++, Tcl/Tk.	
o How extensible and/or customizable is the tool. source code given tool implemented in a public-domain language not extensible by user other:	
3. TOOL FEATURES AND UTILITIES	
o Tool supports the following (check all that apply): _X GUI for GCCS	

	the library is validated
	The extent of the library is (speaking from the point of view of a potential user): not very comprehensive reasonably comprehensive quite comprehensive
_x	Editing and document preparation tools _textual user interface for VPL
	Cross-referencing Browsing
	Increments tracing Incremental development across multiple sessions Change control and version management Consistency checking
	Other: graphical compiler for generating Facile code (similar
	to Standard ML and CCS), Java and Ada'95 codegraphical simulators for GCCS_simulator for VPL
_x	ractive/mechanized/automated is the tool. fully automated user guided other:
4. TOOL INF	PUT AND OUTPUT
_X	oorts these models: synchronous asynchronous nixed
o Input to GCCS	the tool. or VPL specification or combination of the two.
Step Step	rom the tool. 1: networks of finite-state processes. 2: model checking and/or code generation.
GCCS:	formal semantics modern programming language constructs (e.g. if-else):
	strong typing modularity
_x	hierarchical design parameterization

	_X communication between processes buffered
	built-in model of computation
	X other: _graphical
	based on CCS
7	/PL:
	_X formal semantics
	modern programming language constructs (e.g. if-else):
	integers of limited size
	arrays and records of integers
	if-then-else
	while-do
	select
	
	strong typing
	modularity
	hierarchical design
	X_ parameterization
	communication between processes buffered
	built-in model of computation
	X other:finite data domain
	A Other
3. TO	OL APPLICATION
o Abst	craction level that the tool can address (check all that apply):
	requirements
	_X design specification
	X implementation
	(code generation)
	test derivation
	RTL
	netlists
	transistor level
	other:
O Had	the tool been integrated with other tools?
O mad	X no
	yes
	with
	do not know
	Company Compan
4. RES	SOURCES
o Reso	ource requirements for the tool:
	UNIX versionSunOS 4.1 or Solaris on Sun SPARC
	Windows version
	Mac version
	Memory:

o Cost, rights and restrictions:
free, no license
X_ free, license required
_X for educational and research use only
A for calculation shares
nominal distribution charge
fee for underlying tool(s)
flat license fee
per user license fee
royalties per use
other:
o User background prerequisites (check all that apply):
_X BS degree
MS degree
Nb degree
Ph.D. degree
knowledge of logic
first-order
high order
x familiarity with a high-level programming language
familiarity with a high level programming language familiarity with process algebra
familiarity with temporal logic
other:
o User's learning curve, if all prerequisites are met:
X one month
A Old Month
two months
less than six months
other
months
o Tool support
_X upgrades/maintenance
Last version produced at this date: _1998_
New version to be released in near future.
manual
on the web
training
listserv
mailing list
mailing list dedicated conference(s)/workshop(s)
human "help line"
book(s)
X journal/conference publications
other:
o Current contact.
concurr@cs.sunysb.edu
http://www.cs.sunysb.edu/~concurr
6. QUESTIONS APPLYING TO MODEL CHECKERS ONLY
o Verification mechanism(s) (check all that apply):
equivalence
X modal logic
linear-time local and global model checker for alteration-free

modal mu-calculus _local model checker for real-time extension for the above logic temporal logic system or process invariants built-in support for checking for: deadlock livelock other:
_X other:strong and weak bisimulation
o Tool supports (check all that apply): LMC (local model checker):
<pre>_X optimization and state reduction mechanism usingon-the-fly execution and partial order reduction_ simulator:</pre>
interactive random
X_ feedback on in what state verification failed _X_ trace leading to the state (if the user chooses so)
7. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
o Degree of proof mechanization. fully mechanized partially mechanized o Support for developing and viewing the proof. o Presentation of proof to the user (e.g., user input or canonical expressions, with or without quantifiers). o Tool supports (check all that apply): automated support for arithmetic reasoning automated support for efficient handling of large propositional expressions automated support for rewriting possible to use lemmas before they are proved possible to state and use axioms without having to prove them new definitions can be introduced and existing definitions modified during proof facilities for editing proofs the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified reasonably easy to reverify a theorem after slight changes to the specification
8. OPEN-ENDED QUESTIONS
o Strengths of this tool. Designed for use by protocol engineers and software developers, for industrial-scale problems. Specification, simulation, verification and code generation of concurren systems modeled as hierarchical networks of finite-state processes.

Sophisticated graphical support for specification and simulation.

Automatic code generation from verified specifications.

o Limitations of this tool.

Finite-state systems.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Main application area is reactive systems, including embedded system software, process control systems, telecommunication protocols, security protocols, and e-commerce protocols.

o Applications that the tool was used for - case studies, examples, success stories.

Posted on http://www.cs.sunysb.edu/~concurr/. Examples:

Specification and verification of: GNU UUCP i-Protocol, E-2C Hawkeye Early

Warning Aircraft Display LAN Protocol, RETHER real-time Ethernet protocol.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal_Methods/

Murphi

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: _X model checker theorem prover mechanized proof assistant other:
o Application domain(s) or class(es) of problems originally intended. Hardware protocol verification, optional extensions for cryptographic protocols.
Early design stages, error finding. o Intended audience. Digital designers.
o Language(s) and/or technique(s) that the tool is based on. Murphi language: collection of guarded rules (condition/action), which are executed repeatedly in an infinite loop (similar to Chandy and Misra's Unity language.)
o Reasoning mechanisms used for the tool. Explicit state space enumeration, depth- or breath- first search; simulation.
o Comparable tools: SMV, Spin, Concurrency Factory, CWB.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation.
o How extensible and/or customizable is the tool. _X source code given _X tool implemented in a public-domain language not extensible by user other:
3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): GUI
Library of standard types, functions, and other constructions the library is validated
The extent of the library is (speaking from the point of view of

	a potential user):
	not very comprehensive
	reasonably comprehensive
	quite comprehensive
Note:	while there is no standard library, a number of types and
	functions that are commonly provided by a library are provided in
	the language, for example, arrays, records, Multiset and
Scalarset.	
	Editing and document preparation tools
	zaroznaj ana accament proportion contra
	Cross-referencing
	Browsing
	Requirements tracing
	Incremental development across multiple sessions
	Change control and version management
	Consistency checking
-	Completeness checking
	Other:
o How inter	ractive/mechanized/automated is the tool.
	fully automated
	user guided
	other:
4 TOOL INE	OUT AND OUTPUT
1. 1002 111	
o Tool gunr	ports these models:
	synchronous
_^	asynchronous
	(interleaving)
	mixed
o Input to	
	description.
	om the tool.
	boolean invariant is violated, error message and error trace.
•	ts if error or assertion statements are reached.
o The langu	age used for input to the tool has (check all that apply):
	formal semantics
_x	modern programming language constructs (e.g. if-else):
	if
	switch
	for
	while
	strong typing
	modularity
	hierarchical design
	narameterization

hffamad
buffered
built-in model of computation
other:
5. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
requirements
X design specification
implementation
implementation test derivation
RTL
netlists
transistor level
other:
o Has the tool been integrated with other tools?
no Y ves
_X yes
withSVC
do not know
do not know
6. RESOURCES
o Resource requirements for the tool:
UNIX versionprecompiled for: INDY IRIX 5.3,
C_{11} C_{12} C_{12} C_{12} C_{13} C
SunSPARC20 SunOS 4.1.3_U1, 4.1.4, 5.4,
SunSPARCserver-1000 SunOS 5.5,
SunSPARCserver-1000 SunOS 5.5, Intel Linux 1.3.48, 2.0.27, 2.0.34, 2.0.36_
SunSPARCserver-1000 SunOS 5.5, Intel Linux 1.3.48, 2.0.27, 2.0.34, 2.0.36_ Windows version
SunSPARCserver-1000 SunOS 5.5, Intel Linux 1.3.48, 2.0.27, 2.0.34, 2.0.36_
SunSPARCserver-1000 SunOS 5.5, Intel Linux 1.3.48, 2.0.27, 2.0.34, 2.0.36_ Windows version
SunSPARCserver-1000 SunOS 5.5, Intel Linux 1.3.48, 2.0.27, 2.0.34, 2.0.36_ Windows version Mac version
SunSPARCserver-1000 SunOS 5.5,

first-order
high order
X familiarity with a high-level programming language
familiarity with process algebra
familiarity with temporal logic
other:
o User's learning curve, if all prerequisites are met:
X one month
two months
less than six months
other:
months
o Tool support
_X upgrades/maintenance
Last version produced at this date: _Murphi 3.1, 1999
X manual
X on the web
training
training listserv
V mailing list
_A Mailing list dedicated conference(s)/workshop(s)
human "helm line"
_x other:bug reports, suggestions to murphi@verify.stanford.edu
ounceag reported, buggereress of samples of the
o Current contact.
http://sprout.stanford.edu/dill/murphi.html
murphi@verify.stanford.edu
marphire to the property of th
7. QUESTIONS APPLYING TO MODEL CHECKERS ONLY

o Verification mechanism(s) (check all that apply):
equivalence
modal logic
temporal logic
X system or process invariants (boolean propositions true
for all states of the system/process)
built-in support for checking for:
X deadlock
livelock
other:error statements
assertion statements
other:
o Tool supports (check all that apply):
X optimization and state reduction mechanism
state reduction using:
symmetry (description has identical elements that
• • • · · · -

	can be permuted consistently without changing
	verification properties)
	reversible rules (condition/action can be executed "in
	reverse")
	repetition constructors (keeping track of how many
processes	<u> </u>
P	are in the same state)
	hash compression algorithms for probabilistic
verification	nash compression argorithms for probabilistic
	zation using:
	probabilistic verification
	state space caching
	parallel Murphi
	using magnetic disk instead of main memory
simulat	or:
i	nteractive
	andom
	k on in what state verification failed
	ace leading to the state
other:	ace reading to the beate
Other:	
8. QUESTIONS ABOUT	THEOREM PROVERS [NASA98]
o Degree of proof	
fully m	echanized
partial	ly mechanized
o Support for deve	loping and viewing the proof.
	proof to the user (e.g., user input or canonical
expressions,	,
with or without	quantifiers)
	heck all that apply):
	ed support for arithmetic reasoning
	ed support for efficient handling of large propositional
express	
automat	ed support for rewriting
possibl	e to use lemmas before they are proved.
possibl	e to state and use axioms without having to prove them.
new def	initions can be introduced and existing definitions
	d during proof
	ies for editing proofs
	ndations (i.e., all axioms, definitions, assumptions,
	of the proof are identified
	bly easy to reverify a theorem after slight changes to
the spe	cification
9. OPEN-ENDED QUES	TIONS
o Strengths of thi	s tool.
	industrial use by non-experts in formal methods.
	and state reduction algorithms and techniques.
o Limitations of t	
	for liveness and fairness properties (e.g. livelock).
NO CHECKTIIG	LOL LIVENCED AND LATINCED DIODELLIES (E.G. LIVELOCK).

No message communication.

Not possible to describe sequential behavior.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Multiprocessor cache coherence protocols. Security protocols. o Applications that the tool was used for - case studies, examples, success stories.

Listed at http://sprout.stanford.edu/dill/murphi.html. Examples: Verification of cache coherence protocols for Sun UltraSparc-1 Verification of cache coherence and link level protocol for Sun's S3.mp multiprocessor

Specification and verification of Sparc V9 TSO, PSO, RMO memory models Cryptographic and security protocols

Verification of a part of "Scalable Coherent Interface" IEEE Std 1596-1992

Proprietary industrial protocols, for Fujitsu, HAL Computer Systems,
HP, IBM, ad others

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1. http://eis.jpl.nasa.gov/quality/Formal_Methods/

SVM Cadence

*************** SMV Cadence Berkeley Labs *************

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification:
X_ model checker
theorem prover
<pre>X mechanized proof assistant (of limited scope and built on top of the model checker) other:</pre>
o Application domain(s) or class(es) of problems originally intended. Hardware verification.
o Intended audience.
General.
o Language(s) and/or technique(s) that the tool is based on.
SMV input language is used to describe a refinement hierarchy (that
is, specifications at multiple levels of abstraction).
Specifications are written in temporal logic, or an HDl-like
equational notation. It is also possible to input models in a
synchronous version of the Verilog HDL. The logic is effectively
a first-order, quantifier free, linear time temporal logic.
o Reasoning mechanisms used for the tool.
Model checking (determines the truth of temporal formulas by exhaustive
state-space exploration).
o Comparable languages/tools.
Spin, the Concurrency Workbench, the Concurrency Factory, VIS,
Mocha, COSPAN, FormalCheck.
This tool is an extension of Carnegie Mellon SMV to support
compositional methods.
Note: SMV is a research vehicle, and is not directly related the
FormalCheck product from Cadence.
Formatcheck product from Cadence.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation.
OBDD-based model checking algorithm, implemented in C language.
Compositional proof methods, also implemented in C.
o How extensible and/or customizable is the tool.
source code given
tool implemented in a public-domain language
X not extensible by the user.
other:

3. TOOL FEATURES AND UTILITIES

o Tool supports the following (check all that apply):
_X GUI
Library of standard types, functions, and other constructions
the library is validated
The extent of the library is (speaking from the point of view
of
a potential user):
not very comprehensive
reasonably comprehensive
quite comprehensive
Note: while there is no standard library, a number of types and
functions that are commonly provided by a library are provided in the
language, for example, bit vectors and binary arithmetic, arrays,
structures. Queues are notably absent, however.
X_ Editing and document preparation tools
Emacs interface
Cross-referencing
X Browsing
Requirements tracing
Incremental development across multiple sessions
Change control and version management
Consistency checking
Completeness checking
X Other:
BDD library (implemented in C) for sequential verification
support for refinement verification
o How interactive/mechanized/automated is the tool.
X fully automated
X user guided
(User guidance is required for refinement verification.)
other:
4. TOOL INPUT AND OUTPUT
1, 1002 20101
o Tool supports these models:
X synchronous
X asynchronous
X mixed
o Input to the tool.
Model in SMV language (a collection of properties expressed
in temporal logic) or Synchronous Verilog (which is then translated
into SMV language). o Output from the tool.
Yes/no answer to posed temporal formulas, counterexample if "no."
Also,
keeps track of the status of proof obligations in compositional
proofs.
o The language used for input to the tool has (check all that apply):
formal semantics

	X modern programming language constructs (e.g. if-else):
	X strong typing
	(typing is used only to enforce symmetry)
	_X modularity
	X hierarchical design
	_X parameterization
	(can describe designs with arbitrary number of components,
etc.)	
	_X communication between processes
	(signals and shared variables)
	buffered
	built-in model of computation
	other:
5. TOOL	APPLICATION
a 7 h a to sa	nation level that the tool can address (check all that apply).
O ADSUL	action level that the tool can address (check all that apply):
	requirements
	X design specification
	X implementation
	test derivation
	_X RTL
	_X netlists
	transistor level
	other:
o Has tl	ne tool been integrated with other tools?
	no
	X yes - please name tool and applications
	with bounded model checker form CMU_
	with
	with
	do not know
6. RESO	URCES
o Resou	rce requirements for the tool:
	UNIX versionIntel 386 Linux, SPARC SunOS, Solaris, HPUX,
MIPS/Ir:	
	Windows versionNT, 95
	Mac version
	Memory:
o Cost,	rights and restrictions:
	free, no license
	X free, license required
	X for educational and research use only
	nominal distribution charge
	fee for underlying tool(s)
	flat license fee
	per user license fee
	per user ricense ree royalties per use
	rojureren per une

	other:
o Use:	r background prerequisites (check all that apply):
	_X BS degree
	MS degree
	Ph.D. degree
	knowledge of logic
	first-order
	high order
	familiarity with a high-level programming language
	familiarity with process algebra _X familiarity with temporal logic other:
	other:
- 77	ala languing gumes, if all propoguigitog are met.
o use:	r's learning curve, if all prerequisites are met:
	one month
	_X two months
	less than six months
	other
	months
o Too	l support
	_X upgrades/maintenance
	Last version produced at this date:1999
	_X manual
	_X on the web
	_X training
	lecture notes and tutorials, on the web
	listserv
	X mailing list
	dedicated conference(s)/workshop(s)
	human "help line" book(s)
	book(s)
	journal/conference publications
	X other: _archive and FAQ, on the web_
	questions and comments to smv-users@cadence.com
	
o Curi	rent contact.
	http://www.cs.cmu.edu/~modelcheck/index.html for older version of SMV
	http://www.cis.ksu.edu/santos/smv-doc/
	http://www-cad.eecs.berkeley.edu/~kenmcmil/
	smv-users@cadence.com
7. OUF	ESTIONS APPLYING TO MODEL CHECKERS ONLY
,	
o Veri	ification mechanism(s) (check all that apply):
O VCI	equivalence
	modal logic
	_X temporal logic
	CTL, LTL
	system or process invariants
	built-in support for checking for:
	deadlock

livelock
other:
other:
o Tool supports (check all that apply): _X optimization and state reduction mechanism using compositional methods: data type reduction,
uninterpreted functions, cone-of-influence reduction, temporal case splitting,
constant propagation, circular compositional proofs, symmetry reductions,
<pre>induction over the natural numbers, refinement verification.</pre>
_X simulator: interactive random
_X feedback on in what state verification failed _X_ trace leading to the state 8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
Note: SMV is not a general purpose theorem prover, but it does provide a special-purpose proof assistant.
o Degree of proof mechanization. fully mechanized _X partially mechanized
o Support for developing and viewing the proof. Graphical browser.
o Presentation of proof to the user (e.g., user input or canonical expressions, with or without quantifiers).
None. o Tool supports (check all that apply):
<pre>_X automated support for arithmetic reasoning</pre>
expressions automated support for rewriting _X possible to use lemmas before they are proved.
X possible to state and use axioms without having to prove themX_ new definitions can be introduced and existing definitions modified during proof _X_ facilities for editing proofs
_X the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified _X reasonably easy to reverify a theorem after slight changes to
the specification 9. OPEN-ENDED QUESTIONS

- o Strengths of this tool.

Combines model checking and compositional proof methods.

This means that, on the one hand, the state explosion

problem can be avoided by decomposition, while on the other hand, model checking can be used to avoid writing detailed invariants by hand.

o Limitations of this tool.

Not user-extensible, in the way that most proof assistants are.

Limited to first-order temporal logic.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Hardware verification.

o Applications that the tool was used for - case studies, examples, success stories.

Verification of the RTL-level implementation of a cache coherence protocol

(SGI), as well as numerous cache coherence protocols at an abstract level.

Verification of standard hardware protocols, e.g. Futurebus+ and PCI local bus protocols.

Numerous applications in low-level hardware verification.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for

Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal_Methods/

SPIN

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification: _X model checker theorem prover mechanized proof assistant other:
o Application domain(s) or class(es) of problems originally intended. Software, distributed systems. o Language(s) and/or technique(s) that the tool is based on.
PROMELA (PROcess MEta LAnguage), a non-deterministic language loosely based on Dijkstra's guarded command language notation, and borrowing the notation for I/O operations from Hoare's CSP
language. o Reasoning mechanisms used for the tool. State space exploration (exhaustive or partial); simulation.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation. ANSI C, on-the-fly checking. o How extensible and/or customizable is the tool. _X source code given _X tool implemented in a public-domain language not extensible by user other:
3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): _X GUI (Xspin) Library of standard types, functions, and other constructions the library is validated
The extent of the library is (speaking from the point of view of a potential user): not very comprehensive reasonably comprehensive quite comprehensive

	Note: while there is no standard library, a number of types and functions that are commonly provided by a library are provided in the language, for example, arrays and queues. Editing and document preparation tools
	Cross-referencing Browsing
	Requirements tracing
	Incremental development across multiple sessions
	Change control and version management
	X Consistency checking
	_X Completeness checking
	Other:
	depository of source code extensions on SPIN web
page	
o How	interactive/mechanized/automated is the tool.
	_X fully automated
	_X user guided
	(simulation option)
	other:
	L INPUT AND OUTPUT supports these models: synchronous _X_ asynchronous (interleaving) mixed
o Inpu	t to the tool.
	Model written in PROMELA (somewhat resembles a C program).
o Outp	ut from the tool.
	Yes/no answer to posed tests;
	trace leading to errors;
	% coverage of state space.
	language used for input to the tool has (check all that apply):
	_X formal semantics X modern programming language constructs (e.g. if-else):
	do
	u-u-u
	strong typing
	modularity
	hierarchical design
	X parameterization
	X communication between processes
	X_ buffered
	X rendezvous X through shared memory
	x through shared memory X built-in model of computation

other:
5. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
X requirements
X design specification
X implementation
_X test derivation
_A Det
RTL
netlists
transistor level
other:
o Has the tool been integrated with other tools?
no
X_ yes
withSCR* toolset for tabular specifications_
with PEP
with
with
do not know
6. RESOURCES
6. RESOURCES
- Deserved manufacture for the tool.
o Resource requirements for the tool:
UNIX versionany standard UNIX, Linux
Windows version95/98, NT
Mac version
Memory:
o Cost, rights and restrictions:
free, no license
X free, license required
X for educational and research use only
nominal distribution charge
fee for underlying tool(s)
flat license fee
per user license fee
royalties per use
other:
**
o User background prerequisites (check all that apply):
_X BS degree
MS degree
Ph.D. degree
knowledge of logic
first-order
high order
X familiarity with a high-level programming language
familiarity with process algebra

	familiarity with temporal logic
	other:
o User's l	earning curve, if all prerequisites are met:
	one month
	two months
	less than six months
	other:
_	months
o Tool sup	
_x	upgrades/maintenance
	Last version produced at this date:Spin 3.3.3, 1999
X	manual
	_X on the web
	training
	listserv
	mailing list
	<pre>mailing list dedicated conference(s)/workshop(s)</pre>
_X	dedicated conference(s)/worksnop(s)
	(annual, international, since 1995)
	human "help line"
_X	book(s)
x	journal publications
	other:regular electronic newsletter
	(mailed out and posted on the web page)
	proceedings of Spin workshops, on the web page
	web page with source code extensions depository
	hug renorts and suggestions to the newsletter
a Charant	bug reports and suggestions, to the newsletter
o Current	contact.
spin	contactlist@research.bell-labs.com (newsletter)
spin	contact.
spin	contact. list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY
spin 7. QUESTION O Verificat	contact. list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply):
spin 7. QUESTION O Verificat	contact. list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply):
spin 7. QUESTION O Verificat	contact. _list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY tion mechanism(s) (check all that apply): equivalence
spin 7. QUESTION O Verificat	contact. _list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic
spin 7. QUESTION O Verificat	contact. _list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic
spin 7. QUESTION O VerificatX	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL
spin 7. QUESTION O VerificatXX	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL system or process invariants
spin 7. QUESTION O VerificatXX	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL system or process invariants other:never claims (Buchi automata)
spin 7. QUESTION O VerificatXX	contact. _list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL system or process invariants other:never claims (Buchi automata) trace can be replayed in simulator to demonstrate
spin 7. QUESTION O VerificatXX	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL system or process invariants other:never claims (Buchi automata)
spin 7. QUESTION O Verificat XXX	contact. _list@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL system or process invariants other:never claims (Buchi automata)
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spin 7. QUESTION O Verificat X X X X X X X X X X X X X X X X X X	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL_ system or process invariants other:never claims (Buchi automata) trace can be replayed in simulator to demonstrate property violation corts (check all that apply): optimization and state reduction mechanism usingpartial order reduction, bit-state hashing (optional), wolper's hash-compact method (optional), storing reachable states with minimized automaton, statement merging, nested depth-first search algorithm
spin 7. QUESTION O Verificat X X X X X X X X X X X X X X X X X X	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic ITL_ system or process invariants other:never claims (Buchi automata)
spin 7. QUESTION O Verificat X X X X X X X X X X X X X X X X X X	contactlist@research.bell-labs.com (newsletter) NS APPLYING TO MODEL CHECKERS ONLY cion mechanism(s) (check all that apply): equivalence modal logic temporal logic LTL_ system or process invariants other:never claims (Buchi automata) trace can be replayed in simulator to demonstrate property violation corts (check all that apply): optimization and state reduction mechanism usingpartial order reduction, bit-state hashing (optional), wolper's hash-compact method (optional), storing reachable states with minimized automaton, statement merging, nested depth-first search algorithm

X guided	
X feedback on in what state verification failed	
X trace leading to the state	
built-in support for checking for:	
Suffer in Support for thecking for: _X_ deadlock	
A deadlock	
X livelock	
X boolean propositions	
X other:LTL formulas (internally converted into never	
claims)	
dynamically growing and shrinking number of	
processes	
semaphores	
semaphores	
O OHECHTONG ADOLER MILEODEM PROLIDEG (MAGAGO)	
8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]	
a Dogway of week week-winstin	
o Degree of proof mechanization.	
fully mechanized	
partially mechanized	
o Support for developing and viewing the proof.	
o Presentation of proof to the user (e.g., user input or canonical	
expressions,	
with or without quantifiers).	
o Tool supports (check all that apply):	
automated support for arithmetic reasoning	_
automated support for efficient handling of large proposition	ìΤ
expressions	
automated support for rewriting	
possible to use lemmas before they are proved.	
possible to state and use axioms without having to prove them	_
new definitions can be introduced and existing definitions	
modified during proof	
facilities for editing proofs the foundations (i.e., all axioms, definitions, assumptions, lemmas) of the proof are identified	
the foundation (i.e., 1)	
the foundations (i.e., all axioms, definitions, assumptions,	
reasonably easy to reverify a theorem after slight changes to	
the specification	
9. OPEN-ENDED QUESTIONS	
o Strengths of this tool.	
Easy to learn by people with some programming experience.	
Optimized for verifying large problem sizes (e.g. bit-state hashing	
on-the-fly checking).	,
Actively contributing user community in more than 40 countries.	
o Limitations of this tool.	
Not efficient to specify large data sets.	
o Estimated possible uses of the tool, such as applications, classes of	
problems, stages of production cycle.	
Develop verified process control systems from requirements to	
implementation.	
Trace logical design errors in distributed systems, such as	
operating systems, railway signaling protocols, data communications	
)
protocols, switching systems, concurrent algorithms. O Applications that the tool was used for - case studies examples	
U AUDITUALIUMS LHAL THE FOOT WAS USED FOR - CASE STUDIES AVAMBLES	

success stories.

Posted throughout Spin News Letters and workshop proceedings, http://netlib.bell-labs.com/netlib/spin/news.

Some examples include: specification, design, verification and implementation of a safe object oriented process control application, verification of Java applications, steam boiler, hardware cache coherence protocols,

NASA's fault tolerant embedded space craft controller, a multi-threaded plan execution programming language of NASA's New Millennium Remote Agent artificial intelligence based spacecraft control system architecture, telecommunications and security protocols,

Dutch mobile sea-level control.

References:

NRL Protocol Analyzer	

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.	
1. GENERAL DESCRIPTION OF THE TOOL	
o Rough classification: _X model checker theorem prover _X mechanized proof assistant other:	
o Application domain(s) or class(es) of problems originally intended. Analysis of cryptographic protocols used to authenticate principals and	
services and distribute keys in a network. Proving properties of security protocols and finding flaws in them. o Intended audience.	
<pre>o Language(s) and/or technique(s) that the tool is based on.</pre>	
o Comparable languages/tools. STeP.	
2. TOOL IMPLEMENTATION	
o Underlying mechanism of the tool's implementation. Prolog. o How extensible and/or customizable is the tool. source code given x tool implemented in a public-domain language other:	
3. TOOL FEATURES AND UTILITIES	

o Tool supports the following (check all that apply):

	GUI
	Library of standard types, functions, and other constructions the library is validated
	The extent of the library is (speaking from the point of view of a potential user): not very comprehensive reasonably comprehensive quite comprehensive
	Editing and document preparation tools
	
	Cross-referencing Browsing Requirements tracing Incremental development across multiple sessions
	Change control and version management Consistency checking Completeness checking Other:
_x	ractive/mechanized/automated is the tool. fully automated user guided possible to switch between automated and manual mode other:
4. TOOL IN	PUT AND OUTPUT
	oorts these models: synchronous asynchronous mixed
of lo	ription of state in terms of words known by intruder and values ocal state variables.
Compl paths	om the tool. ete description of all reachable states and non-redundant that may precede the specified state. failed/passed.
o The langu	age used for input to the tool has (check all that apply): formal semantics

modern programming language constructs (e.g. if-else):
strong typing
modularity
hierarchical design
_X parameterization
communication between processes
buffered
built-in model of computation
other:
3. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
requirements
X_ design specification
implementation
test derivation
RTL
netlists transistor level
other:
o Has the tool been integrated with other tools?
no
yes
with
with
withwith
do not know
Interface for a requirements language.
Interface for high-level security language CAPSL.
4. RESOURCES
a Dagayyaa wagaiyamanta fan tha taal
o Resource requirements for the tool:
UNIX version
Windows version Mac version
Mac version
o Cost, rights and restrictions:
free, no license
free, license required
receipe required
nominal distribution charge
nominal distribution charge fee for underlying tool(s)
<pre> nominal distribution charge fee for underlying tool(s) free for educational and research use only</pre>

flat license fee	
per user license fee	
royalties per use	
other:	
o User background prerequisites (check all that apply):	
BS degree	
MS degree	
Ph.D. degree	
Ph.D. degree X knowledge of logic	
first-order	
high order	
familiarity with a high-level programming language	
familiarity with process algebra	
familiarity with temporal logic	
other:	
o User's learning curve, if all prerequisites are met:	
one month	
two months	
less than six months	
more than six months	
months	
o Tool support	
X_ upgrades/maintenance	
Last version produced at this date:1999	
manual	
on the web	
training	
listserv	
mailing list	
conference(s)/workshop(s)	
human	
book(s)	
X journal/conference publications	
other:	
o Current contact.	
Catherine Meadows	
Code 5543, Naval Research Laboratory, Washington DC 2037	_
meadows@itd.nrl.navy.mil	~
meadowsertd.nrr.navy.mrr	
http://www.itd.nrl.navy.mil/ITD/5540/projects/crypto.htm	. ד
nttp://www.itd.nri.navy.mri/iid/5540/projects/crypto.ncm	11
C OTTECHTONS ADDIVING TO MODEL CURCUEDS ONLY	
6. QUESTIONS APPLYING TO MODEL CHECKERS ONLY	
Note: we will generate the state evaluation nextical of MDI	
Note: we will consider the state exploration portion of NRL	
Protocol Analyzer as "model checker."	

a milital and a
equivalence
modal logic
temporal logic
system or process invariants
built-in support for checking for:
deadlock
livelock
other:
other:state exploration

o Tool supports (check all that apply):
X optimization and state reduction mechanism
usingnarrowing algorithm
built-in rules for discarding redundant/unreachable
paths and states
_user-generated rules using a database of formal
languages
symbolic simulator:
interactive
random
_X feedback on in what state verification failed
X trace leading to the state
7. QUESTIONS ABOUT THEOREM PROVERS/MECHANIZED PROOF ASSISTANTS [NASA98] Note: we will consider the proof-oriented part of NRL protocol Analyzer as "theorem prover".
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Note: we will consider the proof-oriented part of NRL protocol Analyzer as "theorem prover". o Degree of proof mechanization. fully mechanized partially mechanized o Support for developing and viewing the proof. o Presentation of proof to the user (e.g., user input or canonical expressions, with or without quantifiers). o Tool supports (check all that apply): automated support for arithmetic reasoning automated support for efficient handling of large propositional expressions automated support for rewriting possible to use lemmas before they are proved possible to state and use axioms without having to prove them.
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reasonably easy to reverify a theorem after slight changes to the specification

8. OPEN-ENDED QUESTIONS

- o Capabilities of this tool.
- o Limitations of this tool.
- o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.
- o Applications that the tool was used for case studies, examples, success stories.

Questionnaire for potential users:

o Briefly describe problems that you need solved (in order to help us estimate if those problems can be addressed by formal tools).

- o Have you used formal tools? If yes, for what application? What were the areas of satisfaction? What were the problem areas? What would you like to see in the future?
- o Describe your dream toolkit.
- o What would you consider a "good place" to integrate formal tools in existing or separate toolkits?

Questionnaire for tool makers/integrators:

o If you already produce and/or sell toolkits, would you be interested in integrating formal tools in the toolkit, and why.

o What information do you need in order to be able to integrate formal tools in a toolkit.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal_Methods/

SCR*

******** Software Cost Reduction (SCR*) ***********************************
For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
<pre>o Rough classification:</pre>
2. TOOL IMPLEMENTATION
<pre>O Underlying mechanism of the tool's implementation.</pre>
3. TOOL FEATURES AND UTILITIES
o Tool supports the following (check all that apply): _X GUI Library of standard types, functions, and other constructions the library is validated

The extent of the library is (speaking from the point of view of a potential user): not very comprehensive reasonably comprehensive quite comprehensiveX Editing and document preparation toolsspecification editor for creating requirements specifications
X dependency graph browserX dependency graph browserX Browsing Requirements tracing X Incremental development across multiple sessions Change control and version management X Consistency checking X Completeness checking Other:
simulator, with visual front ends tailored to
X user guided other:
4. TOOL INPUT AND OUTPUT
o Tool supports this kind of models:
synchronous
X_ asynchronous
mixed
o Input to the tool. Tabular SCR specification; asynchronous input from non-deterministic environment.
o Output from the tool.
Specification editor output:
dictionaries with static information (e.g. names of variables user-defined types) tables
Dependency graph browser:
directed graph depicting dependencies among variables
Consistency checker:
syntax and type errors, missing cases, variable name
discrepancies, unwanted nondeterminism, and circular
definitions.
Abstraction derivator:
more abstract model, eliminated irrelevant variables and

o The language used for input to the tool has (check all that apply): X formal semantics
modern programming language constructs (e.g. if-else):
_X strong typing
_X modularity
hierarchical design parameterization
communication between processes
buffered
_X built-in model of computation
other:

3. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
_X requirements
design specification implementation
test derivation
_x RTL
(under current investigation)
netlists transistor level
X other: documentation
levels that can be addressed with Spin and PVS
o Has the tool been integrated with other tools?
no yes
with _Spin model checker
with PVS theorem prover using TAME high-level user interface
with
with do not know
do not know
4. RESOURCES
o Resource requirements for the tool:
UNIX version SunOS
Windows versionfor Oct'99 release Mac version
Memory:
o Cost, rights and restrictions:
free, no license
X for educational and research use only
nominal distribution charge
fee for underlying tool(s)
flat license fee
per user license fee

	<pre>_X system or process invariants _X_ other:never claims (Buchi automata)trace can be replayed in simulator to demonstrate</pre>
	modal logic _X temporal logicLTL
o Ver	ification mechanism(s) (check all that apply): equivalence
Note:	this section applies to model checker Spin.
6. QU	ESTIONS APPLYING TO MODEL CHECKERS ONLY
	http://www.chacs.itd.nrl.navy.mil/SCR
	kirby@itd.nrl.navy.mil labaw@itd.nrl.navy.mil
	Naval Research Laboratory, Code 5546, Washington DC 20375
o Cur	rent contact.
	_X journal/conference publications other:
	Istserv mailing list dedicated conference(s)/workshop(s) human "help line" book(s)
	dedicated conference(s)/workshop(s)
	mailing list
	_X training
	on the web
	Last version produced at this date: _1998 X manual
	_X upgrades/maintenance
о Тоо	months 1 support
	other
	less than six months
	X one month two months
o Use	r's learning curve, if all prerequisites are met:
	other:
	familiarity with temporal logic
	familiarity with a high-level programming language familiarity with process algebra
	high order
	first-order
	Ph.D. degree knowledge of logic
	MS degree
o Use	r background prerequisites (check all that apply): X BS degree
	royalties per use other:
	rovalties per use

property violation
o Tool supports (check all that apply): _X_ optimization and state reduction mechanism usingpartial order reduction, bit-state hashing (optional), Wolper's hash-compact method (optional), storing reachable states with minimized automaton, statement merging, nested depth-first search algorithm
_X interactive _X random _X guided _X feedback on in what state verification failed X trace leading to the state built-in support for checking for: _X deadlock _X livelock _X_ boolean propositions _X_ other:LTL formulas (internally converted into never
claims) dynamically growing and shrinking number of processessemaphores
7. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
Note: this section applies to mechanized proof assistant PVS, with TAME interface and SCR validity checker.
<pre>o Degree of proof mechanization: fully mechanized _X partially mechanized (although finite state verification and the proof of many straightforward results are fully automatic. There is also a batch mode in which proofs may be easily rerun, and a facility for defining proof strategies to automate proofs. o Support for developing and viewing the proof: Tcl/Tk interface to display proof trees and theory hierarchies. Proofs yield scripts that may be edited, attached to additional</pre>
formulas, and rerun. Proofs may also be checkpointed, providing rapid access to parts of a proof the user wishes to examine or adjust.
o Presentation of proof to the user (e.g., user input or canonical
expressions with or without quantifiers): Proofs are presented in a sequent-style representation. PVS takes care to assure that the initial proof goal transparently reproduces the formula input by the user. Quantification is retained; implicit universal quantification in the user's specification is made explicit o Tool supports (check all that apply):
X automated support for arithmetic reasoning X automated support for efficient handling of large propositional

	expressions
_X	automated support for rewriting
_x	possible to use lemmas before they are proved.
	possible to state and use axioms without having to prove them
_x	new definitions can be introduced and existing definitions
	modified during proof
X	facilities for editing proofs
_x	the foundations (i.e., all axioms, definitions, assumptions,
	lemmas) of the proof are identified
Х	reasonably easy to reverify a theorem after slight changes to
	the specification
Х	other:
	integration with CTL model checking
	ground evaluator (providing "run" speeds comparable to
	imperative programs)
	proof strategies
	proof storage, replay, and checkpointing
	graphical display of proof trees, theory hierarchies, and
	prover commands
	proof chain analysis
	proof and theory status reporting

8. OPEN-ENDED QUESTIONS

o Capabilities of this tool.

Mathematically founded tool for non-specialists in formal methods. Well-developed user interface.

o Limitations of this tool.

Flat structure of specifications.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Requirements specification, specification, verification, documentation. o Applications that the tool was used for - case studies, examples, success stories.

Listed in

http://www.itd.nrl.navy.mil/ITD/5540/personnel/heitmeyer.html. Avionics systems, telephone networks, nuclear power plants, etc.: English-language requirements for NASA International Space Station. Requirements specification for flight guidance system.

Specification and verification of contractor-developed: Weapons Control Panel, and a cryptographic system.

References:

[NASA98] for

NASA, "Formal Methods Specification and Verification Guidebook

Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal Methods/

Tatami			

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.			
1. GENERAL DESCRIPTION OF THE TOOL			
o Rough classification: model checker theorem prover _X mechanized proof assistant _X other: _integrated suite of tools: Kumo, web-based proof			
2. TOOL IMPLEMENTATION			
o Underlying mechanism of the tool's implementation. Java 1.2, OBJ3. o How extensible and/or customizable is the tool source code given _X tool implemented in a public-domain language			
other:			

3. TOOL FEATURES AND UTILITIES

o Tool supports the following (check all that apply): _X GUI	
Library of standard types, functions, and other constructions the library is validated	
The extent of the library is (speaking from the point of view a potential user): not very comprehensive	of
reasonably comprehensive quite comprehensive	
_X Editing and document preparation tools	
Cross-referencing	
_X Browsing Requirements tracing	
<pre>_X Incremental development across multiple sessions _X Change control and version management Consistency checking</pre>	
Completeness checking X Other: executing proof scores on a remote server	
o How interactive/mechanized/automated is the tool. _X fully automated _X user guided other:	
4. TOOL INPUT AND OUTPUT	
o Tool supports these models: synchronous asynchronous mixed	
o Input to the tool. Specification in BOBJ; prof script with execution commands in Duck language.	
o Output from the tool. Proof results. Kumo generates web pages with documentation based on user input.	
o The language used for input to the tool has (check all that apply): _X formal semantics X modern programming language constructs (e.g. if-else):	

o Res	with
o Res	with do not know SOURCES Ource requirements for the tool: UNIX version Windows version Mac version t, rights and restrictions: free, no license free, license required nominal distribution charge fee for underlying tool(s) free for educational and research use only flat license fee per user license fee royalties per use
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	with do not know SOURCES ource requirements for the tool: UNIX version
	with with do not know SOURCES
6. RE	with with do not know
6 77	with with do not know
	withwith
	with
	with CafeORJ environment
	no X_ yes - please name tool and applications
o наs	the tool been integrated with other tools?
	other:
	transistor level
	RTL netlists
	test derivation
	X implementation
	X design specification
	X requirements
o Abs	traction level that the tool can address (check all that apply):
5. 10	V# *** * ## V*** # V&1
5 TO	OL APPLICATION
	other:
	built-in model of computation
	buffered
	_X parameterization communication between processes
	nierarcnical design
	<pre>strong typing X_ modularity hierarchical design</pre>

0		X B	round prerequisites (check all that apply): S degree S degree
		Pl	h.D. degree
		kı	nowledge of logic
			first-order
		-	high order
		f	amiliarity with a high-level programming language
			amiliarity with a high-level programming language amiliarity with process algebra amiliarity with temporal logic
		f:	amiliarity with temporal logic
			ther:
			Addition to the second
O	User	's lear	rning curve, if all prerequisites are met:
_		01	ne month
		ti	wo months
).	ess than six months
		of	ther
			months
_	Tool	suppor	months
U			ogrades/maintenance
			Last version produced at this date: _1999
			dast version produced at this date: _1999
		X ma	anua l
			X_ on the web
		+ 1	raining
			istserv
			ailing list
	•		for CafeOBJ
			edicated conference(s)/workshop(s)
	•		for CafeOBJ
			uman "help line"
	•	X bo	
	•		For OBJ3
			ournal/conference publications
			cher:
0	Curr	ent cor	ntact.
•			/www-cse.ucsd.edu/groups/tatami/
	•		www.coc.aca, groups, casame,
7.	QUE	STIONS	APPLYING TO MODEL CHECKERS ONLY
0	Veri	ficatio	on mechanism(s) (check all that apply):
_		ec	quivalence
	-	mc	odal logic
	-	te	emporal logic
	-	sv	stem or process invariants
	-	bu	vstem or process invariants wilt-in support for checking for:
	-		deadlock
		-	livelock
		-	other:
			ther:
0	Tററി ¯		ts (check all that apply):
_	1001		otimization and state reduction mechanism
	-		

using
symbolic simulator:
interactive
random
feedback on in what state verification failed
trace leading to the state
8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
o Degree of proof mechanization.
fully mechanized
partially mechanized
o Support for developing and viewing the proof.
Web-based.
o Presentation of proof to the user (e.g., user input or canonical
expressions,
with or without quantifiers).
- M1
o Tool supports (check all that apply):
automated support for arithmetic reasoning
automated support for efficient handling of large propositional
expressions
automated support for rewriting
possible to use lemmas before they are proved.
possible to state and use axioms without having to prove them.
new definitions can be introduced and existing definitions
modified during proof
facilities for editing proofs
the foundations (i.e., all axioms, definitions, assumptions,
lemmas) of the proof are identified
reasonably easy to reverify a theorem after slight changes to
the specification

9. OPEN-ENDED QUESTIONS

o Capabilities of this tool.

Ease of use, user interface and system operation designed for software engineers who are not experts in formal methods.

Will be possible to use various proof checkers other than Kumo.

o Limitations of this tool.

Kumo is not a powerful proof assistant like HOL or PVS.

o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.

Cooperative web-based software system design and validation. o Applications that the tool was used for - case studies, examples, success stories.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal_Methods/

What Tool Makers Need for Tool Integration (1 received response)

Questionnaire for tool makers/integrators:

o If you already produce and/or sell toolkits, would you be interested in integrating formal tools in the toolkit, and why.

Integration is happening. Need a spectrum of tools for any kind of useful system.

o What information do you need in order to be able to integrate formal tools in a toolkit.

API, sockets main link into Z/EVES.

Appendix B: Formal Methods Term Taxonomy

Formal Methods Term Taxonomy

Background

Mature life-cycle process, in the context of system engineering, consists of: requirements definition, system design, high-level design, low-level design, implementation, testing (unit testing, component testing, and system testing), user support, and maintenance.

Model is a system of definitions, assumptions and equations, set up to represent and discuss physical phenomena and systems. In the context of mathematical logic, a model is an implementation, I, of a set of well-formed formulas of a formal language such that each member of the set is true in I.

Axiom is a mathematical formula that can assert arbitrary properties over arbitrary (new or existing) entities.

Definition, is an axiom that introduces a new symbol and gives its value or meaning as a function of previously existing symbols.

Theorem is a logical formula derived from axioms using inference rules.

Method, in the context in an engineering discipline, describes a way in which a process is to be conducted. In the context of system engineering, a method consists of: 1) underlying model of development; 2) a language or languages; 3) defined ordered steps; and 4) guidance for supplying them in a coherent manner.

Proof is a chain of reasoning using rules of inference and a set of axioms that leads to conclusion, i.e. it is derivation of a theorem.

Step-wise refinement, in the context of system engineering, is the process of deriving level i+1 of the process cycle from level i, and refining level i based on level i+1, in systematic fashion through all cycles of life-cycle.

Taxonomy

Abstraction is the process of simplifying and ignoring irrelevant details and focusing, distilling, and generalizing what remains. In formal methods, abstraction is a tool for eliminating

distracting detail, avoiding premature commitment to implementation choices, and focusing on the essence of the problem at hand.

Breadth-first search is a search that generates first all the immediate neighbors of a state, then all the next neighbors, and so on.

Completeness is a property defined as presence of all possible cases.

Consistency is a property defined as lack of conflicting cases.

Explicit model checking is a type of model checking in which the system to be analyzed is represented by enumerating its states and transitions. State exploration is performed over individual states. The term "model checking" usually implies explicit model checking.

Formal analysis is mathematically-based analysis.

Formal method is a mathematically-based technique for describing system components, properties and/or behavior. Formal methods are different than traditional engineering mathematics in the sense that they are used for describing digital systems, such as hardware and software, using logic and discrete mathematics. A formal method has an underlying theoretical model against which a description can be verified. It consists of a notation (i.e. formal specification language) and some form of deductive apparatus (i.e. proof system).

Formal methods may be applied at varying levels of rigor or formalization. Listed in order of increasing formality and effort, a suggestive guide to levels of rigor includes:

- 1. Use of notations and concepts derived from logic and discrete mathematics to develop more precise requirements statements and specifications. Proof, if any, is informal.
- 2. Use of formalized specification languages with mechanized support tools ranging from syntax checkers and prettyprinters to typecheckers.
- 3. Use of fully formal specification languages with rigorous semantics and correspondingly formal proof methods that support theorem proving and model checking.

Formal proof is a complete and mathematically based argument for the validity of a statement about a system description. A proof proceeds in a series of steps, each of which draws conclusions from a set of assumptions. Justification for each step is derived from a small set of rules which state what conclusions can be reasonably drawn from assumptions. Such justification eliminates ambiguity and subjectivity from the argument. Formal proofs may be prepared manually or, preferably, with the assistance of a formal methods tool.

Formal specification is a description of a planned or existing process, entity and/or system, written in a formal language. It is a concise and unambiguous description of the behavior and/or properties of the process/entity/system, and can be written at various levels of abstraction and formalization. It can be used for requirements, system design, high-level design, and low-level

design specification, as well as test derivation. The most formal specifications are written in languages with well-defined semantics that support formal deduction and allow the consequences of the specification to be calculated through proof of putative theorems.

Formal (specification) language is a mathematically based language, and has a formal syntax and semantics.

- Formal languages can be broadly classified as model-oriented, property-oriented, or a combination of both. Model-oriented languages explicitly model system behavior. Property-oriented language describe properties of the system.
- Formal languages can also be classified as sequential or concurrent, if they are used to specify
 sequential or concurrent systems, respectively. For example, process algebras are model-oriented
 languages which describe the behavior of concurrent systems by describing their algebra of
 communicating processes.
- Formal languages can be executable, and can have tool support.
- Programming languages are formal languages, but are not considered appropriate for use in formal specifications because of: insufficient abstraction ability (e.g. in "true" formal languages, types do not have to be directly implementable); often there is a lack of complete formal semantics.

Formal (methods) tool is a program that implements some aspect of formal analysis, thus providing mechanized, computer assisted support for formal analysis. Like formal methods, formal methods tools can be formalized to various levels of rigor, from syntax checkers to theorem provers.

Formal validation is a type of formal analysis in which an implementation is tested in execution to demonstrate that it satisfies its requirements specification. Informally, it is proving that the requirements are right, (i.e. we are building the desired system).

Formal verification is a form of formal analysis in which each level of development is proven to satisfy the requirements of its superior level, (i.e. formal specification satisfies the corresponding formal requirements specification, and implementation satisfies the corresponding formal specification). Informally, it is proving that a system is built to its requirements.

Formalization is the application of a certain level of mathematical rigor; or the act of formalizing an informal process, system or entity by making it more mathematically rigorous. In the context of using formal languages and tools, levels of formalization are (in increasing order):

- 1. Use of mathematical concepts and notation, informal analysis (if any), absence of mechanized assistance.
- 2. Use of formalized specification language with some mechanical support.
- 3. Use of formal specification language with comprehensive mechanized environment, which includes mechanized proof assistant/theorem prover and/or model checker.

Mechanized proof assistant is a formal tool that implements theorem proving in an interactive way, requiring the user to guide the proof steps.

Model checking is a type of formal analysis that relies on building a (usually finite) model of a system and checking that a desired property holds in that model. The verification task is to demonstrate that the system is a model that satisfies the putative property. The specification should be syntactically and semantically correct. The check is performed as an exhaustive or partial state space search, often breadth-first. Model checking is based on a verification algorithm and thus requires no assistance from the user, i.e. it is "automatic."

Model checker is a formal tool that implements model checking. Model checkers usually rely on various algorithms, such as bit-state hashing or symmetry, to reduce state space search, and/or in the case of very large systems could provide an option to perform nearly exhaustive state space search.

Theorem proving is a type of formal analysis in which a proof of a property is performed over a specification. Both the specification and its properties are expressed as formulas in some kind of mathematical logic. The verification task is to show that the formal specification of the system implies the formal statement of a putative system property. The specification should be syntactically and semantically correct.

Theorem prover is a formal tool that implements theorem proving in an automated way, not requiring user assistance.

Parser is a formal tool that checks syntactic consistency.

Requirements specification is a specification describing essential, necessary or desired attributes of a system or system components.

Rule of inference is a rule in mathematical logic that defines the reasoning that determines when a conclusion may be drawn from a set of premises. In a formal system, the rules of inference should guarantee that *if* the premises are true, *then* the conclusion is also true.

Specification animators (or **emulators**) are executable programs which reinterpret a formal specification into a high-level dynamically executable form. Specification animations are not formal in a strict sense, but support the formal requirements and design verification process by providing analysts with an early view of the high-level dynamic behavior of the requirements.

Symbolic execution is execution which does not require parameters to have known values, (i.e., allows parameters in symbolic form).

Symbolic model checking is an approach to model checking in which the system to be analyzed is described by equations or logical formulas. For example, a form of symbolic model checking uses the state reduction technique to analyze sets of states, represented as Boolean formulas,

instead of individual states. For illustration, let us consider the state in which V is set to 0. All states that have V set to 0 are marked, and all states that can reach the marked states in one step are marked. This procedure is repeated until no new states can be marked. This set of states is then analyzed.

Symbolic simulation is a form of simulation that allows input parameters to be supplied in symbolic form, (e.g. as variables or functions).

Traceability of requirements is a property which means that system-level requirements are traceable to identifiable (functional) subsystems, components, or interfaces.

Typechecking is a form of formal analysis that detects semantic inconsistencies and anomalies, ensuring that entities must match their declaration and be combined only with other entities of the same or compatible type.

Typechecker is a formal tool that implements typechecking.

Unparser (or pretty-printer) is a tool that translates internal representations into display, and outputs formatted text. Usually used at the specification level.

Questionnaire

Tools Makers/Users

For this particular tool, please answer the following questions grouped based on: general description of the tool, tool implementation, tool features and utilities, applications and resources.
1. GENERAL DESCRIPTION OF THE TOOL
o Rough classification:
model checker
theorem prover
mechanized proof assistant other:
other: other: o Application domain(s) or class(es) of problems originally intended.
o Intended audience.
o Language(s) and/or technique(s) that the tool is based on.
o Reasoning mechanisms used for the tool.
o Comparable languages/tools.
2. TOOL IMPLEMENTATION
o Underlying mechanism of the tool's implementation.
o How extensible and/or customizable is the tool.
source code given tool implemented in a public-domain language
not extensible by user
other:

3. TOOL	FEATURES AND UTILITIES
o Tool - -	supports the following (check all that apply): GUI Library of standard types, functions, and other constructions the library is validated
	The extent of the library is (speaking from the point of view of a potential user): not very comprehensive reasonably comprehensive quite comprehensive
	Editing and document preparation tools
	
 	Cross-referencing Browsing Requirements tracing Incremental development across multiple sessions Change control and version management Consistency checking Completeness checking Other:
o How i	nteractive/mechanized/automated is the tool fully automated user guided other:

4. TOOL INPUT AND OUTPUT
o Tool supports these models: synchronous asynchronous mixed
o Input to the tool.
o Output from the tool.
o The language used for input to the tool has (check all that apply): formal semantics modern programming language constructs (e.g. if-else): strong typing modularity hierarchical design parameterization communication between processes buffered built-in model of computation other:
5. TOOL APPLICATION
o Abstraction level that the tool can address (check all that apply):
o Has the tool been integrated with other tools? no yes - please name tool and applications with with with with

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0	Resource	requirements for the tool:
	UNIX	version
	Windo	ows version
	Mac v	version
		EY:
0	Cost, rig	hts and restrictions:
		free, no license
		free, license required
		for educational and research use only nominal distribution charge
		nominal distribution charge
		fee for underlying tool(s)
		flat license fee
		per user license fee
		royalties per use
		other:
0		ground prerequisites (check all that apply):
		BS degree
		MS degree
		Ph.D. degree
		knowledge of logic
		first-order
		nian Order
		familiarity with a high-level programming language
		familiarity with process algebra
		familiarity with a high-level programming language familiarity with process algebra familiarity with temporal logic
		other:
0	User's le	earning curve, if all prerequisites are met:
		one month
		two months
		less than six months
		more than six months
		months
	_	
0	Tool supp	
		upgrades/maintenance
		Last version produced at this date:
		manual
		on the web
		training
		listserv
		mailing list
		dedicated conference(s)/workshop(s)
		human "help line" book(s)
		DOOK(S)
		journal/conference publications
		other:
_	Current a	ont a ct

7. QUESTIONS APPLYING TO MODEL CHECKERS ONLY
<pre>o Verification mechanism(s) (check all that apply):</pre>
other:
o Tool supports (check all that apply): optimization and state reduction mechanism using simulator: interactive random simbolic feedback on in what state verification failed trace leading to the state
other:

8. QUESTIONS ABOUT THEOREM PROVERS [NASA98]
o Degree of proof mechanization. fully mechanized partially mechanized o Support for developing and viewing the proof.
o Presentation of proof to the user (e.g., user input or canonical expressions, with or without quantifiers).
o Tool supports (check all that apply): automated support for arithmetic reasoningautomated support for efficient handling of large propositional expressionsautomated support for rewritingpossible to use lemmas before they are provedpossible to state and use axioms without having to prove themnew definitions can be introduced and existing definitionsmodified during prooffacilities for editing proofsthe foundations (i.e., all axioms, definitions, assumptions,lemmas) of the proof are identifiedreasonably easy to reverify a theorem after slight changes tothe specification

- 9. OPEN-ENDED QUESTIONS
- o Capabilities of this tool.
- o Limitations of this tool.
- o Estimated possible uses of the tool, such as applications, classes of problems, stages of production cycle.
- o Applications that the tool was used for case studies, examples, success stories.

References:

[NASA98] NASA, "Formal Methods Specification and Verification Guidebook for Software and Computer Systems", vol.1.

http://eis.jpl.nasa.gov/quality/Formal_Methods/

Questionnaire

Potential Users

- Briefly describe problems that you need solved (in order to help us estimate if those problems can be addressed by formal tools)
- ♦ Have you used formal tools? If yes,
 - ◆ For what application?
 - What were the areas of satisfaction?
 - ♦ What were the problem areas?
 - ♦ What would you like to see in the future?
- Describe your dream toolkit.
- ♦ What would you consider a "good place" to integrate formal tools in existing or separate toolkits?

Questionnaire

Tools Makers/Integrators

- ♦ If you already produce and/or sell toolkits, would you be interested in integrating formal tools into the toolkit, and why?
- What information do you need in order to be able to integrate formal tools into a toolkit?

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